

# Approximate Commutative Algebra – an Impossible Concept ?

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# Algebra and **Approximation** a Fundamental Antithesis

## Algebra

discrete world

symbols, structures

no neighborhoods

equal  $\langle \Rightarrow \rangle$

not equal

satisfied  $\langle \Rightarrow \rangle$

not satisfied

true  $\langle \Rightarrow \rangle$  false

discrete

elements

equations

relations

mappings

## Analysis

continuous world

spaces, functions

neighborhoods

equal ... close

... distant

sat. ... nearly sat.

... not satisfied

nearly true ... false

continuous

**Approximation** can only be considered

in an **analytic** context, hence

(Commutative) Algebra and **Approximation**

are **antithetic** concepts

But many problems in Commutative Algebra are posed over number fields which carry a

**natural topology**

e.g. real numbers or complex numbers

Here,  $1.41421 \approx \sqrt{2}$  may be  
an obvious **approximation**

This suggests an

**embedding**

of **commutative algebra** into **analysis**  
via the coefficient field

Successful Prototype for this Embedding :

Linear Algebra over  $\mathbb{R}^n$  or  $\mathbb{C}^n$

becomes

**Analytic Geometry**

becomes

**Numerical Linear Algebra**

**Embedding** into **Analysis** permits

- introduction of **approximation**
- handling of data with **limited accuracy**
- use of **analytic tools** for problem solving
- use of **floating-point** in computation

**Embedding** requires

- **reconsideration** of all concepts
- **redefinition** of many concepts
- introduction of various **new concepts**
- **disposal** of some classical concepts

**Embedding** provides

- new **insights**
- means to **avoid ill-conditioning**
- **faster** algorithms
- **safer** algorithms

These benefits arise even for **exact** computation with **exact** data !

**Analytic** treatment of an **algebraic** problem  
has been introduced by

**C. F. Gauss**

Solution of systems of linear equations  
by **iterative approximation**

**A Revolutionary Idea :**

Have an **algebraic** problem  
whose (exact) solution may be obtained by  
a *finite number* of (exact) *rational operations*

**Instead** take an approximate solution and  
**improve it iteratively** and

**Stop** when **accuracy is sufficient**  
e.g. relative to *data accuracy*

Developed and used by Gauss  
in a large-scale surveying project  
Communicated (by letter) to another surveyor  
on Dec. 26, 1823

Numerous details of the **numerical** computation  
are touched upon in this communication :

- computation of a good initial approximation
- work with *corrections* not with values
- meaningful accuracy during iteration
- details of choosing the iteration sequence
- stopping criterion

**Today , in Scientific Computing ,**

- “all” linear systems solved in **floating-point**
- “almost all” linear systems solved **iteratively**

where “solved”  $\equiv$  solved **approximately**

**Is this **meaningful** and **achievable**  
also for **polynomial problems** ?**

# Fundamental Idea for the **Analytic Treatment** of Quantitative **Algebraic** Problems

“quantitative”: Besides by their structure, problem and results are characterized by **numbers**

Problem : “data”       $a \in \mathcal{A} \subseteq \mathbb{C}^M$

Result : “solution”       $z \in \mathcal{Z} \subseteq \mathbb{C}^m$

## Model of Quantitative Problem :

Given: a **mapping**  $F : \mathcal{A} \rightarrow \mathcal{Z}$  (implicitly!)

Sought:  $F(a)$  for specified  $a \in \text{dom}F \subset \mathcal{A}$

$F$  is the **exact** “data→result mapping”

**structural** aspects of problem are contained in  $F$

definition of  $F$  may include values like 1, 0, etc.

*must include structural sparsity*

**Note:** a coefficient 0 may signify

- an absent term (*sparsity*)

- a tiny value (*approximate data*)

## Analytic Properties of $F : \mathcal{A} \rightarrow \mathcal{Z}$

$$\dim \mathcal{A} =: M, \quad \dim \mathcal{Z} =: m$$

$F$  is **well-posed** in a neighborhood  $N(a)$  of  $a$  if

domain  $F \supset N(a)$ , *open* in  $\mathcal{A}$

$m \leq M$  and  $F$  is *surjective* in  $N(a)$

$F$  is *continuous* in  $N(a)$

otherwise it is **ill-posed**

$F$  is **well-conditioned** in  $N(a)$  if

$F$  is *well-posed* in  $N(a)$

$F$  is *differentiable* in  $N(a)$

$\|F'\|$  is “small”

$F$  is **ill-conditioned** if  $\|F'\|$  “large” in  $N(a)$

Ill-conditioning implies that the solution  $z$  is *highly sensitive* to certain changes in  $a$



## Ill-posed Algebraic Problems

**Type 1 :**  $\dim (\text{domain of } F) =: d < M$

Results only defined for *data on a manifold*  $\mathcal{S}$   
of dimension  $d$  in  $\mathcal{A}$  (i.e. only for special data)

**Type 2 :**  $\dim (\text{image of } F) =: \bar{d} < m$

All exact *results* lie on a *manifold*  $\bar{\mathcal{S}}$   
of dimension  $\bar{d}$  in  $\mathcal{Z}$

(i.e. **approximate** solution is **not** a proper result)

**Type 3 :**  $F$  is *discontinuous* in  $N(a)$

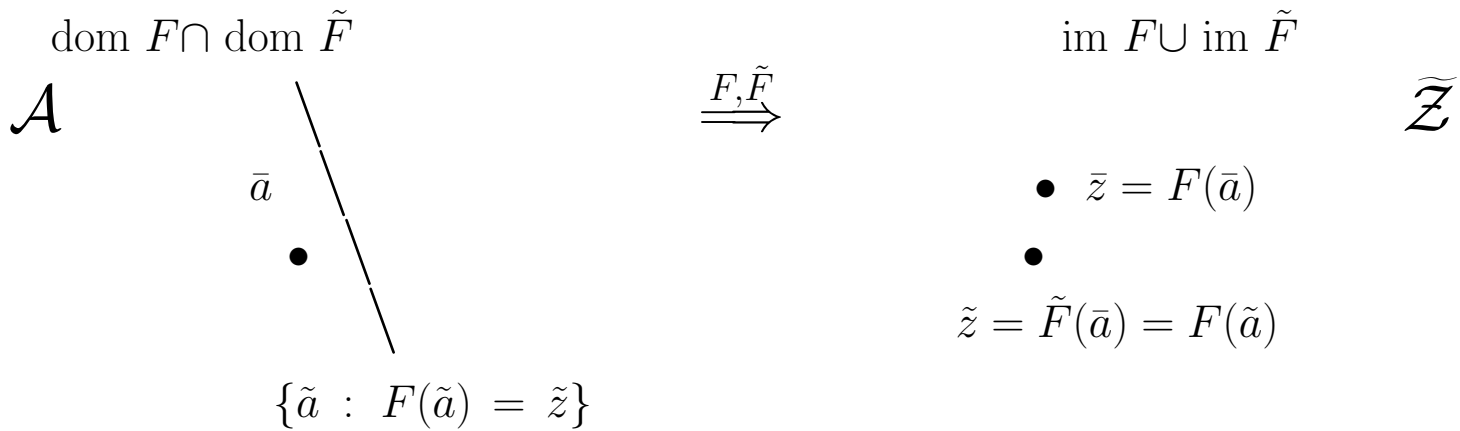
i.e. results “jump” at nearby data

Many classes of interesting quantitative  
algebraic problems are **ill-posed**

Meaningful *approaches* to **ill-posed** problems are  
an important task of Approximate Algebra

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# Approximate Commutative Algebra



## Tasks :

- design **approximate map**  $\tilde{F}$   
(algorithm, implementation)
- evaluate  $\tilde{F}$  for  $\bar{a} \Rightarrow \tilde{z}$   
(run code)
- find  $\tilde{a}$  **closest to**  $\bar{a}$ , with  $F(\tilde{a}) = \tilde{z}$   
(**backward error**)
- correct  $\tilde{z}$  towards  $\bar{z}$   
(**successive approximation**)
- find **analytic** properties of  $F$  around  $\bar{a}$   
(**wellposedness, condition, nearsingularity**)
- re-define  $F$  and/or  $\tilde{F}$  if appropriate  
(respect **analytic** properties of  $F$ )

## Topologies in $\mathcal{A}$ and $\mathcal{Z}$

Formally, we have

$$\text{data space } \mathcal{A} \equiv \mathbb{C}^M$$

$$\text{result space } \mathcal{Z} \equiv \mathbb{C}^m$$

Note: **Structural** data are **not** components of  $\mathcal{A}$

May use some **weighted** norm to define distance:

$$\text{2-norm : } \quad \| \tilde{a} - a \|_{e,2} := \left( \sum_{\nu} \left( \frac{|\tilde{\alpha}_{\nu} - \alpha_{\nu}|}{\varepsilon_{\nu}} \right)^2 \right)^{1/2}$$

$$\text{max-norm : } \quad \| \tilde{a} - a \|_{e,\infty} := \max_{\nu} \frac{|\tilde{\alpha}_{\nu} - \alpha_{\nu}|}{\varepsilon_{\nu}}$$

Weights  $\varepsilon_{\nu}$  should relate to **accuracy levels** of  $\alpha_{\nu}$

e.g.  $\bar{\alpha}_{\nu}$  specified to 4 decimals  $\Rightarrow \varepsilon_{\nu} = 10^{-4}$

$$e := (\varepsilon_1, \dots, \varepsilon_M)$$

Different data components  $\alpha_{\nu}$

may have different **tolerances**  $\varepsilon_{\nu}$

## Empirical Data

data from measurements, observations etc. have

limited accuracy

specified data  $\bar{a} \in \mathcal{A} \subset \mathbb{R}^M$  or  $\mathbb{C}^M$

tolerance  $e \in \mathbb{R}_+^M$

validity parameter  $\delta \in \mathbb{R}_+$

family of sets of **equivalent data** ( $\delta > 0$ ) :

$$N_\delta(\bar{a}, e) := \{\tilde{a} \in \mathcal{A} : \|\tilde{a} - \bar{a}\|_e \leq \delta\}$$

**validity scale** (*continuous gradual* transition!)

$\delta$ : 0 ... 1 ... 3 ... 10 ... 30 ...  
valid    probably    possibly    probably    invalid  
          valid        valid        invalid

**Empirical polynomial** in  $\mathbb{C}^s$  :

$$p(x; a) := \sum_{\nu} \alpha_{\nu} x^{j_{\nu}} \quad \text{with} \quad x^{j_{\nu}} := x_1^{j_{\nu 1}} \dots x_s^{j_{\nu s}}$$

where *some* (or all) coefficients  $\alpha_{\nu}$  are **empirical**

## Valid Instances of Empirical Data

$\tilde{a} \in \mathcal{A}$  is a **valid instance** for  $(\bar{a}, e)$  if

$$\tilde{a} \in N_\delta(\bar{a}, e) \quad \text{with } \delta = \text{“O(1)”}$$

i.e.  $\tilde{a}$  is *indistinguishable* from  $\bar{a}$   
on the specified accuracy level,

## Valid Results for Empirical Data

**computational** result  $\tilde{z} = \tilde{F}(\bar{a}) \in \mathcal{Z}$

*interpreted* as **exact** result for **modified** data  $\tilde{a}$

Equivalent-data manifold  $\mathcal{M}(\tilde{z})$  :

$$\mathcal{M}(\tilde{z}) := \{ \tilde{a} \in \mathcal{A} : \tilde{a} \xrightarrow{F} \tilde{z} \} \subset \mathcal{A}$$

E.g. Polynomial zeros :

$$p(x; \bar{a}) := \bar{\alpha}_0 + \bar{\alpha}_1 x + \dots + \bar{\alpha}_{M-1} x^{M-1} + x^M$$

$$\mathcal{M}(\tilde{z}) := \{ \tilde{a} \in \mathcal{A} : p(\tilde{z}; \tilde{a}) = 0 \} \quad \text{linear (M-1)-manifold}$$

$$\tilde{z} \text{ } \delta\text{-valid} \Leftrightarrow N_\delta(\bar{a}, e) \cap \mathcal{M}(\tilde{z}) \neq \emptyset$$

Find *minimal*  $\delta$  :  $\tilde{z}$  is **valid** if  $\delta_{\min} = \text{O(1)}$

# Validity Checking

Have to **compute**

$$\delta_{\min}(\tilde{z}) := \min (\|\tilde{a} - \bar{a}\|_e, \tilde{a} \in \mathcal{M}(\tilde{z}))$$

For a *linear* manifold  $\mathcal{M}(\tilde{z})$ , this is a

- linear *least squares* problem with a *2-norm*
- linear *optimization* problem with a *max-norm*

E.g. Polynomial zeros :

Find *nearest* coefficient set  $\tilde{a}_{\min}$  with  $p(\tilde{z}; \tilde{a}_{\min}) = 0$

$$\delta_{\min}(\tilde{z}) := \|\tilde{a}_{\min} - \bar{a}\|_e$$

Validity checking permits the following

## Algorithmic Scheme

- compute **approximate result**  $\tilde{z}$  *somehow*  
(approximate algorithm, floating-point)
- check validity of  $\tilde{z}$

if **not valid** : refine  $\tilde{z}$  , go to check

else : **accept**  $\tilde{z}$  as **valid result**

## Valid Result Sets

Problem with **empirical** data  $(\bar{a}; e)$  :

$$Z_\delta(\bar{a}; e) := \{ \tilde{z} \in \mathcal{Z} : \tilde{z} \text{ a } \delta\text{-valid result} \}$$

**size** of  $Z_\delta(\bar{a}; e)$  indicates **condition** of result  $F(\bar{a})$

“Pseudospectra” introduced for matrix eigenvalues in 90’s

Valid result sets for *distinct* results **not** meaningful when these results are to be used *jointly*.

Must be *jointly validated* if jointly used ( $k \leq M$ ):

$$\mathcal{M}(\{\tilde{z}_1, \dots, \tilde{z}_k\}) := \mathcal{M}(\tilde{z}_1) \cap \dots \cap \mathcal{M}(\tilde{z}_k)$$

E.g. *Several* polynomial zeros  $\tilde{z}_\kappa$ ,  $\kappa = 1(1)k$  :

$$\mathcal{M}(\{\tilde{z}_1, \dots, \tilde{z}_k\}) := \{\tilde{a} \in \mathcal{A} : p(\tilde{z}_\kappa; \tilde{a}) = 0, \kappa = 1(1)k\}$$

$\{\tilde{z}_1, \dots, \tilde{z}_k\}$  *simultaneously*  $\delta$ -valid iff

$$N_\delta(\bar{a}, e) \cap \mathcal{M}(\{\tilde{z}_1, \dots, \tilde{z}_k\}) \neq \emptyset$$

$$Z_\delta[\{\tilde{z}_1, \dots, \tilde{z}_k\}] \not\subseteq Z_\delta[\tilde{z}_1] \cup \dots \cup Z_\delta[\tilde{z}_k]$$

$M$  pol. zeros:  $\mathcal{M}(\{\tilde{z}_1, \dots, \tilde{z}_M\}) = \{\tilde{a}^*\}$ , interpolating pol.

## Validity for Ill-Posed Problems

Type 1:  $\mathcal{S} := \text{domain}(F) \subset \mathcal{A}$ ,  $\dim \mathcal{S} =: d < M$

$$\mathcal{M}(\tilde{z}) := \{ \tilde{a} \in \mathcal{S} : F(\tilde{a}) = \tilde{z} \} \subset \mathcal{S} \subset \mathcal{A}$$

E.g., *Multiple* zero of  $p(x; \bar{a})$ : discriminant  $(p, p') = 0$   
now, refinement becomes (linearized) *optimization*

Type 2:  $\bar{\mathcal{S}} := \text{image}(F) \subset \mathcal{Z}$ ,  $\dim \bar{\mathcal{S}} =: \bar{d} < m$

$\tilde{z} \notin \bar{\mathcal{S}}$  is *not a result*,  $\mathcal{M}(\tilde{z})$  is *not defined*

Must check validity of  $\tilde{z}$  in further use

Refine  $\tilde{z}$  towards  $F(\bar{a})$  on  $\bar{\mathcal{S}}$

E.g., numerically computed Groebner basis polynomials do *not* satisfy syzygies, may refine towards “exact” syzygies

Type 3:  $F$  discontinuous, e.g. discrete

$\tilde{z}$  valid if  $\mathcal{M}(\tilde{z}) \supset N_\delta(\bar{a}, \mathbf{e})$ ,  $\delta = O(1)$

else: Check validity of “more degenerate” results

E.g., rank of **empirical** matrix = *lowest* valid rank



## Basic Principles

interpret quantitative problem as a

data  $\rightarrow$  result map  $F : \mathcal{A} \rightarrow \mathcal{Z}$

Design approximate map  $\tilde{F}$

Evaluate  $\tilde{F}$  for specified data

obtain approximate result  $\tilde{z}$

Compute backward error of  $\tilde{z}$

**accept** or refine  $\tilde{z}$

**Note :**

Want exact results (except for round-off)

for valid algebraic problems

i.e. for data very close to specified ones

## Conclusions

Approximate Commutative Algebra  
is meaningful and feasible  
in the context of  
quantitative problems  
with coefficients with a natural topology  
it is indispensable for  
empirical problems