

Recent Progress for Computing Gröbner Bases

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New Criterion (G., Volny and Wang 2011)

$g_1, \dots, g_m \in R = \mathbb{F}[x_1, \dots, x_n]$: any given polynomials
 $e_1, \dots, e_m \in R^m$: the unit vectors

Consider the R -submodule of $R^m \times R$:

$$\begin{aligned} M &= \langle (e_1, g_1), \dots, (e_m, g_m) \rangle_R \\ &= \{(\mathbf{u}, v) \in R^m \times R : \mathbf{u}g^t = v\} \end{aligned}$$

where $\mathbf{u}g^t = u_1g_1 + \dots + u_mg_m$.

Theorem

Suppose G is a subset of M containing $(e_1, g_1), \dots, (e_m, g_m)$. For any term order on R and any compatible term order on R^m , the following are equivalent:

- (a) G is a strong Gröbner basis for M ,
- (b)
- (c) every J -pair of G is covered by G .

Condition (c)

Let $G = \{(\mathbf{u}_1, v_1), (\mathbf{u}_2, v_2), \dots, (\mathbf{u}_r, v_r)\} \subset R^m \times R$. We say

- a pair $p = (\mathbf{u}, v) \in R^m \times R$ with $v \neq 0$ is **covered** by G if there is a pair $p_i = (\mathbf{u}_i, v_i) \in G$ and a monomial $t \in R$ so that

$$\text{lm}(\mathbf{u}) = t \text{lm}(\mathbf{u}_i), \quad \text{and} \quad t \text{lm}(v_i) \prec \text{lm}(v).$$

In this case, we say p_i covers p . **Transitive relation**

- a pair $(\mathbf{u}, 0) \in R^m \times R$ is **covered** by G if there is a pair $(\mathbf{u}_i, 0) \in G$ and a monomial $t \in R$ so that

$$\text{lm}(\mathbf{u}) = t \text{lm}(\mathbf{u}_i).$$

Remark. The condition (c) corresponds to F5 rewritten rules and is used in F5, Arri and Perry (2011) and in many recent papers.

Condition (c): Example

Let $R = \mathbb{F}[x, y, z]$ under the graded lex order with $x > y > z$. In $R^3 \times R$, suppose G contains the following pairs:

$$p_4 = (yze_1 + \cdots, 0), \quad p_5 = (xze_2 + \cdots, 0), \quad p_6 = (ze_1 + \cdots, y + \cdots).$$

Then

$$\begin{array}{lll} (xz^2e_2 + \cdots, 0) & \text{is covered by} & p_5, \\ (y^2ze_1 + \cdots, z + \cdots) & \text{is covered by} & p_4, \\ (xze_1 + \cdots, x^2 + \cdots) & \text{is covered by} & p_6, \\ (xze_1 + \cdots, y^2 + \cdots) & \text{is not covered by} & p_6 \end{array}$$

as $xp_6 = (xze_1 + \cdots, xy + \cdots)$ and $y^2 \prec xy \prec x^2$. (In fact, the last pair is not covered by G .)

Definition

For any **term order** in $R = \mathbb{F}[x_1, \dots, x_n]$, a subset $G = \{g_1, \dots, g_r\}$ of an ideal $\mathbf{I} \subset R$ is called a **Gröbner basis** (GB) for \mathbf{I} if every $f \in \mathbf{I}$ is **top-reducible** by G , that is, there exists some $g \in G$ such that $\text{lm}(g)$ divides $\text{lm}(f)$.

The corresponding reduction is

$$f := f - ctg$$

where $t = \text{lm}(f)/\text{lm}(g)$ is a monomial and $c = \text{lc}(f)/\text{lc}(g) \in \mathbb{F}$.

Definition

Let $f, g \in R$. The **S-polynomial** of f and g is defined to be

$$S(f, g) = t_1 f - ct_2 g$$

where $c = \text{lc}(f)/\text{lc}(g)$ and

$$t_1 = \frac{\text{lcm}(\text{lm}(f), \text{lm}(g))}{\text{lm}(f)}, t_2 = \frac{\text{lcm}(\text{lm}(f), \text{lm}(g))}{\text{lm}(g)}.$$

- For example, $f = 4x^3y^4 + \dots$, and $g = 5x^4yz^2 + \dots$,

$$S(f, g) = (5xz^2)f - (4y^3)g = xz^2(4x^3y^4 + \dots) - \frac{4}{5}y^3(5x^4yz^2 + \dots).$$

Buchberger's Criterion (1965)

Theorem

A subset $G = \{g_1, \dots, g_r\}$ of an ideal $\mathfrak{I} \subseteq R$ is a Gröbner basis for \mathfrak{I} iff every S -polynomial of G can be top-reduced to zero by G .

Detecting useless S -polynomials

- Buchberger (1979): If $\gcd(\text{lm}(g_i), \text{lm}(g_j)) = 1$ then $S(g_i, g_j)$ can be top-reduced to 0 by G .
- Lazard (1983), Möller, Mora and Traverso (1992):

syzygies \longleftrightarrow "reduction to 0".

For $(g_1, \dots, g_m) \in R^m$, its syzygy module is defined as

$$H = \{\mathbf{u} = (u_1, \dots, u_m) \in R^m : u_1g_1 + \dots + u_mg_m = 0\}.$$

- Faugère (F5, 2002): Introduces signatures and uses principal syzygies to detect useless S -polynomials.

Recent papers

- Bardet (PhD Thesis, 2006), Stegers (2006), Gash (PhD thesis, 2008), Eder and Perry (2009), Sun and Wang (2009),
- Hashemi and Ars (2010), Sun and Wang (2010), G., Guan and Volny (2010), Zobnin (2010),
- G., Volny and Wang (2010/2011), Volny (PhD Thesis, 2011),
- Huang (2010), Eder and Perry (2010),
- Arri and Perry (2011), Eder and Perry (2011), Eder, Gash, Perry (2011), Sun and Wang (2011), Bigatti, Caboara and Robbiano (2011),
- Roune and Stillman (2012), Galkin (2012), Sun and Wang (2012),
- Eder (2013), Eder and Roune (2013), Gerdt and Hashime (2013), Pan, Hu and Wang (2013), Sun and Wang (2013),
- Simões (PhD thesis, 2013), Sun (2013).
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Let $R = \mathbb{F}[x_1, \dots, x_n]$. A monomial in R is written as

$$x^\alpha = x_1^{a_1} \cdot x_2^{a_2} \cdots x_n^{a_n}.$$

A term in R^m is of the form $x^\alpha \mathbf{e}_i$, $1 \leq i \leq m$.

Fix any term order \prec_1 on R and any term order \prec_2 on R^m . We say they are **compatible** if, for each $1 \leq i \leq m$,

$$x^\alpha \prec_1 x^\beta \text{ iff } x^\alpha \mathbf{e}_i \prec_2 x^\beta \mathbf{e}_i.$$

Faugère (2002): For any polynomial $v \in I = \langle g_1, \dots, g_m \rangle$, **the signature** of v is

$$\min\{\text{lm}(\mathbf{u}) : \mathbf{u} = (u_1, \dots, u_m) \in R^m \text{ and } u_1g_1 + \dots + u_mg_m = v\}.$$

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Definition

For any $(\mathbf{u}, v) \in R^m \times R$, we call $\text{lm}(\mathbf{u})$ **the signature** of (\mathbf{u}, v) .

This is much easier to use in practice!

Top-reductions

Let $p_1 = (\mathbf{u}_1, v_1), p_2 = (\mathbf{u}_2, v_2) \in R^m \times R$ be any two pairs.
When v_2 is nonzero, we say p_1 is top-reducible by p_2 if

- (i) v_1 is nonzero and $\text{lm}(v_2)$ divides $\text{lm}(v_1)$; and
- (ii) $\text{lm}(t\mathbf{u}_2) \preceq \text{lm}(\mathbf{u}_1)$ where $t = \text{lm}(v_1)/\text{lm}(v_2)$.

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The corresponding **top-reduction** is then

$$p_1 - ctp_2 = (\mathbf{u}_1 - ct\mathbf{u}_2, v_1 - ctv_2), \quad (1)$$

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where $c = \text{lc}(v_1)/\text{lc}(v_2)$.

Such a top-reduction is called **regular**, if

$$\text{lm}(\mathbf{u}_1 - ct\mathbf{u}_2) = \text{lm}(\mathbf{u}_1),$$

and **super** otherwise.

When $v_2 = 0$, we say that p_1 is **top-reducible** by $(\mathbf{u}_2, 0)$ if \mathbf{u}_1 and \mathbf{u}_2 are both nonzero and $\text{lm}(\mathbf{u}_2)$ divides $\text{lm}(\mathbf{u}_1)$.

When $v_2 = 0$, we say that p_1 is **top-reducible** by $(\mathbf{u}_2, 0)$ if \mathbf{u}_1 and \mathbf{u}_2 are both nonzero and $\text{lm}(\mathbf{u}_2)$ divides $\text{lm}(\mathbf{u}_1)$.

So the signature of p_1 remains the same under a regular top-reduction but becomes smaller under a super top-reduction.

Recall that, for any $g_1, g_2, \dots, g_m \in R = \mathbb{F}[x_1, \dots, x_n]$,

$$\begin{aligned} M &= \langle (\mathbf{e}_1, g_1), \dots, (\mathbf{e}_m, g_m) \rangle_R \\ &= \{(u_1, \dots, u_m, v) \in R^m \times R : v = u_1 g_1 + \dots + u_m g_m\}. \end{aligned}$$

Definition

A subset G of M is called a **Strong Gröbner basis for M** if every pair in M is top-reducible by some pair in G .

Strong GB \implies GB for I and GB for syzygies

Suppose that $G = \{(\mathbf{u}_1, v_1), \dots, (\mathbf{u}_k, v_k)\} \subset R^m \times R$ is a strong Gröbner basis for M . Then

- 1 a Gröbner basis for the syzygy module of $\mathbf{g} = (g_1, \dots, g_m)$ is

$$\mathbf{G}_0 = \{\mathbf{u}_i : v_i = 0, 1 \leq i \leq k\},$$

- 2 and a Gröbner basis for $I = \langle g_1, \dots, g_m \rangle$ is

$$G_1 = \{v_i : 1 \leq i \leq k\}.$$

Let $p_1 = (\mathbf{u}_1, v_1), p_2 = (\mathbf{u}_2, v_2) \in R^m \times R$ be any two pairs.
We form a J-pair only if v_1 **and** v_2 **are both nonzero**.

Recall the S-polynomial of v_1 and v_2 is $t_1 v_1 - ct_2 v_2$ where $c = \text{lc}(v_1)/\text{lc}(v_2)$, and

$$t = \text{lcm}(\text{lm}(v_1), \text{lm}(v_2)), \quad t_1 = \frac{t}{\text{lm}(v_1)}, \quad t_2 = \frac{t}{\text{lm}(v_2)}.$$

For pairs, we have

$$t_1 p_1 - ct_2 p_2 = (t_1 \mathbf{u}_1 - ct_2 \mathbf{u}_2, t_1 v_1 - ct_2 v_2).$$

Suppose

$$T = \max(t_1 \text{lm}(\mathbf{u}_1), t_2 \text{lm}(\mathbf{u}_2)) = t_i \text{lm}(\mathbf{u}_i)$$

where $i \in \{1, 2\}$.

Definition

If $\text{lm}(t_1 \mathbf{u}_1 - ct_2 \mathbf{u}_2) = T$ then

- T is called the **J-signature** of p_1 and p_2 , and
- $t_i p_i$ is called the **J-pair** of p_1 and p_2 , denoted as $J(p_1, p_2)$.

Examples

$R = \mathbb{F}[x, y, z]$ under the graded lex order with $x \prec y \prec z$, $R^3 \times R$ with TOP (term over position) with $\mathbf{e}_3 \prec \mathbf{e}_2 \prec \mathbf{e}_1$.

$$\begin{aligned} p_1 &= (\mathbf{e}_1, 5xy + \cdots), & p_2 &= (\mathbf{e}_2, 3yz + \cdots) \\ p_3 &= (\mathbf{e}_3, 4xz + \cdots), & p_4 &= (x\mathbf{e}_2 + \cdots, x^2 + \cdots) \end{aligned}$$

Some J-pairs:

$$\begin{aligned} (1, 2) : zp_1 &= (z\mathbf{e}_1, 5xyz + \cdots), & J(p_1, p_2) &= zp_1 \\ xp_2 &= (x\mathbf{e}_2, 3xyz + \cdots), & & \text{as } x\mathbf{e}_2 \prec z\mathbf{e}_1 \end{aligned}$$

$$\begin{aligned} (1, 3) : zp_1 &= (z\mathbf{e}_1, 5xyz + \cdots), & J(p_1, p_3) &= zp_1 \\ yp_3 &= (y\mathbf{e}_3, 4xyz + \cdots), & & \text{as } y\mathbf{e}_3 \prec z\mathbf{e}_1. \end{aligned}$$

Theorem (G, Volny and Wang)

Suppose G is a subset of M that contains $(\mathbf{e}_1, g_1), \dots, (\mathbf{e}_m, g_m)$.
Then the following are equivalent:

- (a) G is a strong Gröbner basis for M ,
- (b) every J -pair of G is eventually super top-reducible by G ,
- (c) every J -pair of G is covered by G .

Corollary

Any J -pair satisfying (c) should be discarded.

Special cases:

- Syzygy rule
- F5 rewritten rule

Example. Let $R = \mathbb{F}[x, y, z]$ under the graded lex order with $x > y > z$. In $R^3 \times R$, suppose G contains the following pairs:

$$p_4 = (yz\mathbf{e}_1 + \cdots, 0), \quad p_5 = (xz\mathbf{e}_2 + \cdots, 0), \quad p_6 = (z\mathbf{e}_1 + \cdots, y + \cdots).$$

Then

$(xz^2\mathbf{e}_2 + \cdots, 0)$	is covered by	p_5 , syzygy rule
$(y^2z\mathbf{e}_1 + \cdots, z + \cdots)$	is covered by	p_4 , syzygy rule
$(xz\mathbf{e}_1 + \cdots, x^2 + \cdots)$	is covered by	p_6 , F5 Rewritten rule
$(xz\mathbf{e}_1 + \cdots, y^2 + \cdots)$	is not covered by	p_6

as $xp_6 = (xz\mathbf{e}_1 + \cdots, xy + \cdots)$ and $y^2 \prec xy \prec x^2$. (In fact, the last pair is not covered by G .)

- **Store only the signature $\text{lm}(\mathbf{u})$, not the whole vector \mathbf{u} .**
This gives Gröbner basis for \mathbf{I} and the minimal leading term of the syzygy module.
- **Use principal syzygies.** Any two pairs $p_1 = (\mathbf{u}_1, v_1)$ and $p_2 = (\mathbf{u}_2, v_2)$ give a principal syzygy:

$$v_2 p_1 - v_1 p_2 = (\mathbf{u}, 0).$$

- **Delete** every J-pair that is covered by a pair in G or H (recorded syzygies), or by another J-pair.
- The J-pairs can be processed in any order.

Theorem

The GVW algorithm terminates in finitely many steps if the term orders \prec_1 on R and \prec_2 on R^m are compatible.

- Huang (2010) proved the case when pairs are processed by increasing order. He also gives a nice example to show that the algorithm may not terminate if the signature order and the polynomial order are not compatible.
- There are several flawed proofs for F5: e.g. Hashemi and Ars (2010), Arri and Perry (2011).

Theorem

If the J -pairs are processed in increasing order, then one gets a minimal strong Gröbner basis.

This means that the algorithm does not perform reduction to any extra J -pairs other than the ones in the minimal basis.

Specific Term Orders

Let \prec be some term order on R . We extend \prec to R^m as follows.

- (POT) The first is called position over term ordering (POT). We say that $x^\alpha \mathbf{E}_i \prec x^\beta \mathbf{E}_j$ if $i < j$ or $i = j$ and $x^\alpha \prec x^\beta$.
- (TOP) The second is the term over position ordering (TOP). We say that $x^\alpha \mathbf{E}_i \prec x^\beta \mathbf{E}_j$ if $x^\alpha \prec x^\beta$ or $x^\alpha = x^\beta$ and $i < j$.

- (g1) Next is the \mathbf{g} -weighted degree followed by TOP. We say that $x^\alpha \mathbf{E}_i \prec x^\beta \mathbf{E}_j$ if $\deg(x^\alpha g_i) < \deg(x^\beta g_j)$ or $\deg(x^\alpha g_i) = \deg(x^\beta g_j)$ and $x^\alpha \mathbf{E}_i \prec_{top} x^\beta \mathbf{E}_j$ where \deg is for total degree.
- (g2) Finally, we have \mathbf{g} -weighted \prec followed by POT. We say that $x^\alpha \mathbf{E}_i \prec x^\beta \mathbf{E}_j$ if $\text{lm}(x^\alpha g_i) \prec \text{lm}(x^\beta g_j)$ or $\text{lm}(x^\alpha g_i) = \text{lm}(x^\beta g_j)$ and $x^\alpha \mathbf{E}_i \prec_{pot} x^\beta \mathbf{E}_j$. Called Schreier order.

Under the POT order, our algorithm corresponds with the G2V algorithm.

Under the ordering $\mathbf{g1}$, our algorithm is related to the XL algorithm but much faster.

GVW algorithm under different term orders

Test Case (# gen)	POT (G2V)	TOP	g1	g2
Katsura5 (22)	4.32	0.91	1	0.65
Katsura6 (41)	14.21	5.76	6.29	3.75
Katsura7 (74)	169.63	33.1	34.66	19.9
Katsura8 (143)	1994.86	214.91	224.18	137.39
Schrans-Troost (128)	2106.48	81.86	85.2	95.62
F633 (76)	71.74	42.8	44.78	36.64
Cyclic 6 (99)	111.81	7539.49	7296.54	128.51
Cyclic 7 (443)	44078.6	-	-	24237.8

Table : Runtime in seconds using Singular 3110 on an Intel Core 2 Quad 2.66 GHz processor

Recent works and Open problems

- Yao Sun (August 2013): View GVW algorithm as **GB conversion** via the MMM algorithm of Marinari, Möller, and Mora (1992) or the FGLM algorithm of Faugère, Gianni, Lazard, and Mora (1993).
- Bruno Simões (PhD Thesis, April 2013): Use Hilbert functions.

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- Improve on bounds for the degree D so that

$$\{u_1g_1 + \cdots + u_mg_m : u_i \in \mathbb{F}[x_1, \dots, x_n], \deg(u_i) \leq D\}$$

contains a Gröbner basis. This is closely related to the Castelnuovo-Mumford regularity (assuming g_i 's are homogeneous).

Thank you!