

Johann Radon Institute for Computational and Applied Mathematics Austrian Academy of Sciences (ÖAW)



Walsh functions, scrambled (0, m, s)-nets, and negative covariance: applying symbolic computation to quasi-Monte Carlo integration

J. Wiart, E. Wong

RICAM-Report 2020-24

Walsh functions, scrambled (0, m, s)-nets, and negative covariance: applying symbolic computation to quasi-Monte Carlo integration

Jaspar Wiart^{a,1,2}, Elaine Wong^{a,1,*}

^aAltenberger Straße 69, 4040 Linz, Austria

Abstract

We investigate base b Walsh functions for which the variance of the integral estimator based on a scrambled (0, m, s)-net in base b is less than or equal to that of the Monte-Carlo estimator based on the same number of points. First we compute the Walsh decomposition for the joint probability density function of two distinct points randomly chosen from a scrambled (t, m, s)-net in base bin terms of certain counting numbers and simplify it in the special case t is zero. Using this, we obtain an expression for the covariance of the integral estimator in terms of the Walsh coefficients of the function. Finally, we prove that the covariance of the integral estimator is negative when the Walsh coefficients of the function satisfy a certain decay condition. To do this, we use creative telescoping and recurrence solving algorithms from symbolic computation to find a sign equivalent closed form expression for the covariance term.

Keywords: quasi-Monte Carlo integration, scrambled digital nets, Walsh functions, symbolic computation, creative telescoping, symbolic summation 2010 MSC: 33F10, 65C99, 11K99

^{*}Corresponding author

Email address: elaine.wong@ricam.oeaw.ac.at (Elaine Wong)

 $^{^1\}mathrm{Austrian}$ Academy of Sciences, Johann Radon Institute (RICAM)

²Johannes Kepler University, Linz

1. Introduction

1.1. History of the Problem

Quasi-Monte Carlo methods use low discrepancy point sets and sequences to estimate multidimensional integrals over the unit hypercube:

$$\int_{[0,1)^s} f(x) dx.$$

Roughly speaking, discrepancy measures the overall deviation between the number of points from a point set that are contained in axis-parallel boxes $[\mathbf{x}, \mathbf{y})$ with the number of points that should be in those boxes (i.e. the number of points in the point set divided by the volume of the box); a smaller discrepancy means a better point set. A class of commonly used point sets, (t, m, s)-nets requires that a certain class of boxes (elementary intervals) of a certain size contain exactly the right number of points.

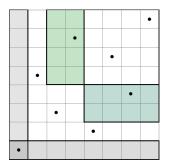


Figure 1: An example of a point set on $[0, 1)^2$ that contains one point in axis-parallel boxes of four different types.

By introducing some randomness into the point sets we can improve the uniform distribution of points and gain access to probabilistic error estimates. It is known that the convergence of the variance of an estimator based on a scrambled (t, m, s)-net (as the number of points increases) is faster than that of the independent and uniformly selected random points in a Monte Carlo (MC) estimator [1, Theoerm 3.9]. Thus, for any given function in $L^2([0,1)^s)$, that the randomized quasi-Monte Carlo (RQMC) estimator will eventually outperform the MC estimator. However, it is not clear how many points are needed before this good behavior happens and until it happens, the RQMC estimator might be worse for a particular function than that of the MC estimator.

In 2018, Lemieux [6] proposed a framework with which to study when the RQMC estimator does no worse than MC. This is based on the concept of negative dependence. Following this, it was shown in Wiart et al. [14], that scrambled (0, m, s)-nets do no worse for functions that are "quasi-monotone". They did

this by showing that scrambled (0, m, s)-nets are negative lower orthant dependent which allowed them to apply a previous result by Lemieux [6]. This required integrating the joint pdf of a scrambled (0, m, s)-net over closed axisparallel boxes anchored at the origin.

Base b Walsh functions have long been known to work well with digital nets in the same base. In this paper, we apply these functions to the variance decomposition framework in [6]. We do this in Section 3 by computing an explicit formula for the base b Walsh coefficients of the joint pdf of scrambled (0, m, s)nets in base b. This yields a formula for the covariance term in terms of the Walsh coefficients of the function. In Section 4 we discuss decay conditions on the Walsh coefficients of functions and how these relate to the average case. In Section 5 we prove that by assuming a natural decay condition on the Walsh coefficients of a function, the estimator based on a scrambled (0, m, s)-net in base b will do no worse than the Monte Carlo estimator based on the same number of points.

1.2. The Use of Symbolic Computation

One of the aims of this work is to introduce the tools of symbolic computation to quasi-Monte Carlo integration. Symbolic computation is a quickly developing field that is always looking for problems with which to apply the methods. We believe that there is a significant opportunity for such tools to aid in computations that are similar to the ones presented here. In this article, we show that we can reduce our problem (of determining whether or not our estimator does better on average than the purely random case) into a manageable form, from which we could draw our conclusions. In this context, we introduce and explain three different tools ([3],[4],[12]) implemented as packages in the computer algebra system Mathematica to help us with our simplifications. We outline the main ideas now and the exact details are shown in Section 5 with computations in the corresponding Mathematica notebook, freely available for download here: HTTPS://WONGEY.GITHUB.IO/DIGITAL-NETS-WALSH/.

An underlying principle that we use to approach this problem is "guess and then prove." The guessing first involves generating a finite amount of data to find a recurrence that the data satisfies. We can make an ansatz with undetermined coefficients for such a recurrence, and obtain necessary conditions on these coefficients by fitting the data. The corresponding linear system is then solved. This is effectively automated with the GUESS.M [3] package, which takes as input the finite data with an estimate on the coefficient degree bound and order of the recurrence, and outputs a recurrence that fits the data (if there is one).

Initially obtaining a recurrence in this way gave us sufficient motivation to simplify our covariance term into a double sum containing (at worst) sums and products of binomial coefficients, which has the nice property of being holonomic. In our setting, this roughly means that the binomial coefficients satisfy recurrences with polynomial coefficients. From there, the method of creative telescoping [15], which has been implemented in HOLONOMICFUNCTIONS.M [4] was then used to compute a recurrence for our double sum. In summary, guessing gave us a recurrence that is valid on the finite data, but not guaranteed to be valid everywhere. Creative telescoping provided a rigorous (and verifiable) proof that the output recurrence holds for all values in the domain of our parameters. In our situation, the latter yielded a higher order recurrence which we were able to show could be derived from the (lower order) guessed one. Then, together with the comparison of initial values, we can assert that the guessed recurrence is indeed correct. In Section 5, we present the final outcome of this computation, with computational details in the notebook that is published online.

Lastly, we can employ yet another tool from the symbolic computation toolbox SIGMA.M [12] to solve the recurrence. In Section 5, the reader will see that a reasonably nice closed form for the solution of the recurrence was produced, and after a few simplifications, we were able to obtain our main result (Theorem 17).

2. Preliminaries

In this paper, we denote \mathbb{N} as the set of natural numbers including 0 and denote $\tilde{P}_n = {\mathbf{U}_1, \ldots, \mathbf{U}_n} \subseteq [0, 1)^s$ to be an RQMC sampling scheme designed to produce an unbiased estimator for the integral, I(f), of a function $f: [0, 1)^s \to \mathbb{C}$ of the form

$$\hat{I}_n(f) = \frac{1}{n} \sum_{i=1}^n f(\mathbf{U}_i),$$

i.e. we assume the \mathbf{U}_i are uniformly distributed in $[0, 1)^s$ with a possible dependence structure between the $\mathbf{U}'_i s$. Because \tilde{P}_n is a randomized sampling scheme, the variance of the estimator, $\operatorname{Var}(\hat{I}_n(f))$, is of interest. In particular, we will seek to better understand which functions satisfy $\operatorname{Var}(\hat{I}_n(f)) \leq \operatorname{Var}(\hat{I}_{MC,n}(f))$, where $\hat{I}_{MC,n}(f)$ is the Monte Carlo estimator of I(f) based on n points.

Let $\psi(\mathbf{x}, \mathbf{y}) \colon [0, 1)^{2s} \to \mathbb{R}_{\geq 0}$ be the joint probability distribution function (pdf) of two distinct points randomly selected from \tilde{P}_n . Following [6], the RQMC variance decomposes as

$$\operatorname{Var}(\hat{I}_n(f)) = \operatorname{Var}(\hat{I}_{MC,n}(f)) + \frac{n-1}{n} \operatorname{Cov}(f(\mathbf{U}_I), f(\mathbf{U}_J))$$
(1)

where \mathbf{U}_I and \mathbf{U}_J are two distinct randomly selected points from \hat{P}_n (we use I and J rather than i and j to emphasise that the points are randomly selected and view $f(\mathbf{U}_I)$ and $f(\mathbf{U}_J)$ as random variables), and

$$\operatorname{Cov}(f(\mathbf{U}_I), f(\mathbf{U}_J)) = \int_{[0,1)^s} \int_{[0,1)^s} (\psi(\mathbf{x}, \mathbf{y}) - 1) f(\mathbf{x}) f(\mathbf{y}) d\mathbf{x} d\mathbf{y}.$$
 (2)

Clearly we have $\operatorname{Var}(\hat{I}_n(f)) \leq \operatorname{Var}(\hat{I}_{MC,n}(f)) \Leftrightarrow \operatorname{Cov}(f(\mathbf{U}_I), f(\mathbf{U}_J)) \leq 0.$

2.1. Scrambled (t, m, s)-nets

We are mainly interested in a particular kind of RQMC sampling scheme known as scrambled (0, m, s)-nets in base b. They arise by scrambling a (0, m, s)-net which are themselves a special case of (t, m, s)-nets in base b, a class of point sets known to have good distribution properties. These nets are typically constructed by using the digital method introduced by Niederreiter [7] and outlined in detail in [1]. However, we will only need the abstract definition.

Let b be a prime and let $P_n = {\mathbf{V}_1, \ldots, \mathbf{V}_n} \subseteq [0, 1)^s$ be a point set with $n = b^m$ points. For $\mathbf{k} \in \mathbb{N}^s$, we say that P_n is \mathbf{k} -equidistributed in base b if each elementary \mathbf{k} -interval of the form

$$\prod_{j=1}^{s} \left[\frac{a_j}{b^{k_j}}, \frac{a_j+1}{b^{k_j}} \right),$$

where $a_j \in \{0, 1, \ldots, b^{k_j} - 1\}$, contains exactly $b^{m-k_1-\cdots-k_s}$ points. If P_n is k-equi-distributed in base b for all $k \in \mathbb{N}^s$ with $k_1 + \cdots + k_s \leq m-t$, we call P_n a (t, m, s)-net in base b (see Figure 2 for examples). The parameter t measures the quality of the point set with smaller values of t being better. The case t = 0 is the best possible. However, a (0, m, s)-net in base b only exists if $b \geq s - 1$ (see [8]*Corollary 4.21).

The goal of scrambling a (t, m, s)-net in base b is to create a randomized version $\tilde{P}_n = {\mathbf{U}_1, \ldots, \mathbf{U}_n} \subseteq [0, 1)^s$ of P_n in such a way that each point \mathbf{U}_i uniformly distributed in the unit hypercube while preserving equidistribution properties. As in [14], a *scrambled* (t, m, s)-net in base b is a (t, m, s)-net that has been digitally scrambled in base b (see Definition 3).

2.2. The Joint PDF of Scrambled (t, m, s)-Nets

We now set some important notation that will be useful when working with the joint pdf of scrambled (t, m, s)-nets.

Definition 1. For $x, y \in [0,1)$, we let $\gamma_b(x, y)$ denote the exact number of initial common digits shared by x and y in their base b expansions, chosen to be finite whenever possible, i.e. the smallest i such that $\lfloor b^i x \rfloor = \lfloor b^i y \rfloor$ but $\lfloor b^{i+1}x \rfloor \neq \lfloor b^{i+1}y \rfloor$. For $\mathbf{x}, \mathbf{y} \in [0,1)^s$, we define

$$\gamma_b^s(\mathbf{x}, \mathbf{y}) = (\gamma_b(x_1, y_1), \dots, \gamma_b(x_s, y_s)) \text{ and}$$
$$\gamma_b(\mathbf{x}, \mathbf{y}) = \sum_{j=1}^s \gamma_b(x_j, y_j).$$

Remark. It is possible that a number $x \in [0, 1]$ has two base b representations. When this happens one representation will be finite and the other will terminate

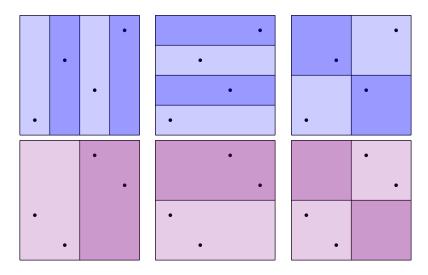


Figure 2: The first row of figures shows the equidistribution properties of a (0, 2, 2)-net in base 2. From left to right we see that the net is (2, 0)-equidistributed, (0, 2)-equidistributed, and (1, 1)-equidistributed in base 2 because each of the corresponding elementary intervals contains one point. The second row of figures shows the equidistribution properties of a (1, 2, 2)-net in base 2. From left to right we see that the net is (1, 0)-equidistributed and (0, 1)-equidistributed in base 2 since each of the corresponding elementary intervals contains two points, but the net is not (1, 1)-equidistributed in base 2.

in an infinite sequence of b-1. For example, $1 = \sum_{i=1}^{\infty} \frac{b-1}{b^i}$. In order for γ_b to be well-defined we must always choose the base b representation of a number to be the finite one whenever possible.

Using $\gamma_b^s(\mathbf{x}, \mathbf{y})$, we define two important classes of sets consisting of pairs of points from $[0, 1)^s$, namely

$$\begin{split} C^s_{\boldsymbol{k}} &= \{ (\mathbf{x}, \mathbf{y}) \in [0, 1)^{2s} : \gamma^s_b(\mathbf{x}, \mathbf{y}) \geq \boldsymbol{k} \} \text{ and } \\ D^s_{\boldsymbol{i}} &= \{ (\mathbf{x}, \mathbf{y}) \in [0, 1)^{2s} : \gamma^s_b(\mathbf{x}, \mathbf{y}) = \boldsymbol{i} \}, \end{split}$$

where the inequality is applied component-wise. When s = 1, we write C_k and D_i . Since

$$C_k = \bigcup_{a=0}^{b^{k+1}-1} \left[\frac{a}{b^k}, \frac{a+1}{b^k}\right]^2$$
 and $D_i = C_i \setminus C_{i+1}$,

 $\operatorname{Vol}(C_k) = b^{-k}$ and $\operatorname{Vol}(D_i) = \frac{b-1}{b^{i+1}}$. Note: $C_k^s = \prod_{j=1}^s C_{k_j}$ and $D_i^s = \prod_{j=1}^s D_{i_j}$. This gives

$$\operatorname{Vol}(C_k^s) = \frac{1}{b^k} \text{ and } \operatorname{Vol}(D_i^s) = \frac{(b-1)^s}{b^{s+i}}.$$

In the above equation we have introduced our convention that, when a letter appears in a formula in both bold and non-bold, the bold letter denotes a vector and the non-bold letter denotes the sum of its coordinates. For example, $i = i_1 + \cdots + i_s$ and $k = k_1 + \cdots + k_s$ for $i, k \in \mathbb{N}^s$.

Definition 2. Let $\tilde{P}_n = {\mathbf{U}_1, \ldots, \mathbf{U}_n} \subseteq [0,1)^s$ be a scrambled net in base $b \ge 2$.

(i) For $\mathbf{k} \in \mathbb{N}^s$, let $M_b(\mathbf{k}; \tilde{P}_n)$ be the number of pairs of distinct points $(\mathbf{U}_l, \mathbf{U}_j)$ in \tilde{P}_n such that $\gamma_b^s(\mathbf{U}_l, \mathbf{U}_j) \geq \mathbf{k}$ (alternatively such that $(\mathbf{U}_l, \mathbf{U}_j) \in C_{\mathbf{k}}^s$). When $\mathbf{k} \in \mathbb{Z}^s$ and \mathbf{k} has a negative component we set

$$M_b(\mathbf{k}; \tilde{P}_n) = M_b(\max(\mathbf{k}; \mathbf{0}), \tilde{P}_n),$$

where the maximum is taken coordinate-wise.

(ii) For $\mathbf{i} \in \mathbb{N}^s$, let $N_b(\mathbf{i}; \hat{P}_n)$ be the number of pairs of distinct points $(\mathbf{U}_l, \mathbf{U}_j)$ in \tilde{P}_n such that $\gamma_b^s(\mathbf{U}_l, \mathbf{U}_j) = \mathbf{i}$ (alternatively such that $(\mathbf{U}_l, \mathbf{U}_j) \in D_{\mathbf{i}}^s$). When $\mathbf{i} \in \mathbb{Z}^s$ and \mathbf{i} has a negative component we set $N_b(\mathbf{i}; \tilde{P}_n) = 0$.

Note that

$$M_b(\mathbf{k}; \tilde{P}_n) = \sum_{\mathbf{k} \le \mathbf{i} \in \mathbb{Z}^s} N_b(\mathbf{i}; \tilde{P}_n)$$
(3)

for all $\mathbf{k} \in \mathbb{Z}^s$ and $M_n(\mathbf{k}; \tilde{P}_n) = b^m(b^{m-k} - 1)$ when \tilde{P}_n is a (0, m, s)-net in base b.

Using this notation we are now able to concisely state our notion of scrambling.

Definition 3. A sampling scheme $\tilde{P}_n = {\mathbf{U}_1, \ldots, \mathbf{U}_n} \subseteq [0, 1)^s$ a base bdigital scramble of $P_n = {\mathbf{V}_1, \ldots, \mathbf{V}_n} \subseteq [0, 1)^s$ if it satisfies the following property: if $(\mathbf{V}_l, \mathbf{V}_j) \in D_i^s$, then $(\mathbf{U}_l, \mathbf{U}_j)$ is uniformly distributed in D_i^s . A scrambled (t, m, s)-net in base b is a (t, m, s)-net that has been digitally scrambled in base b.

One way of realizing such a scramble is Owen's scrambling algorithm [9] (a detailed explanation is given in [1, Section 13.1]).

Theorem 4. (Wiart et al. [14]) Let \tilde{P}_n be a scrambled (t, m, s)-net in base b whose one-dimensional projections are (0, m, 1)-nets. Then the joint pdf $\psi(\mathbf{x}, \mathbf{y})$ of two distinct points randomly chosen from \tilde{P}_n is given by

$$\psi(\mathbf{x}, \mathbf{y}) = \begin{cases} \frac{N_b(i; \tilde{P}_n)}{n(n-1)} \frac{b^{s+i}}{(b-1)^s} & \text{if } i < \infty, \\ 0 & \text{if } i = \infty, \end{cases}$$

where $\mathbf{i} = \gamma_b^s(\mathbf{x}, \mathbf{y})$ and $i = \gamma_b(\mathbf{x}, \mathbf{y})$.

When \tilde{P}_n is a scrambled (0, m, s)-net in base b, the number of pairs of distinct points in \tilde{P}_n that share i initial common digits in their base b expansions can be computed using the formula

$$b^m \sum_{k=0}^{s} (-1)^k \binom{s}{k} \max(b^{m-i-k}, 1)$$

for all $i \in \mathbb{N}^s$ with *i* being the sum of the coordinates of *i* [14]. Since the *t* parameter tells us nothing about the distribution of a (t, m, s)-net P_n on elementary k-intervals where $k_1 + \cdots + k_n > m - t$, we cannot say what the value of the joint pdf will be for a base *b*-digital scramble of P_n without knowing either the points or how the point set was constructed. For this reason, we are unable to obtain a general result for scrambled (t, m, s)-nets when $t \neq 0$.

2.3. Walsh Functions

One aspect of our work in the present paper that differs from the recent work of Wiart et al. [14] is that we take the framework that has already been fixed and investigate the integration of $L^2([0,1)^s)$ functions with Walsh decompositions over the scrambled (0, m, s)-nets. Elementary Walsh functions are piecewise constant and form an orthonormal basis for $L^2([0,1)^s)$. Roughly speaking, they are discrete analogs of sines and cosines. We refer the reader to other sources ([1] [10],[13]) for a complete description of Walsh functions and their properties. We will however, elucidate the properties that we use.

More precisely, for $b \geq 2$, denote ω_b to be the primitive *b*-th root of unity $e^{2\pi i/b}$. Let $l \in \mathbb{N}$ with the (finite) *b*-adic expansion $l = \lambda_0 + \lambda_1 b + \lambda_2 b^2 + \cdots$. Then the *l*-th *b*-adic elementary Walsh function ${}_b \operatorname{wal}_l : \mathbb{R} \to \mathbb{C}$, periodic with period one, is defined

$$_b \operatorname{wal}_l(x) := \omega_b^{\lambda_0 \xi_1 + \lambda_1 \xi_2 + \lambda_2 \xi_3 + \cdots}$$

for $x \in [0, 1)$ with *b*-adic expansion $x = \xi_1 b^{-1} + \xi_2 b^{-2} + \xi_3 b^{-3} + \cdots$. We call $\{b \text{wal}_l : l \in \mathbb{N}\}$ the *b*-adic Walsh function system. For $s \ge 2$ with $x \in [0, 1)^s$ and $l \in \mathbb{N}^s$, we have that

$$_b \operatorname{wal}_l(\boldsymbol{x}) := \prod_{j=1}^s {}_b \operatorname{wal}_{l_j}(x_j).$$

Since the parameter b is symbolically consistent throughout our analysis, we will not include the (pre)-subscript from this point on. Figure 3 gives a pictorial representation of the first few elements of this system in both one and two variables, which illustrate their general behavior on elementary intervals. We now recall some relevant properties of elementary Walsh functions (a more complete list can be found in the sources mentioned above):

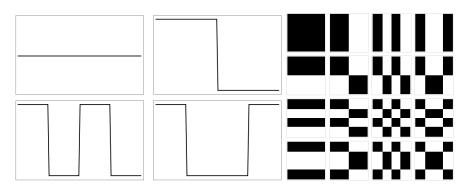


Figure 3: A pictorial example of one-variable (left) and two-variable (right) Walsh functions in base 2 with l = 0, ..., 3 in the one-variable case, $l_1 = 0, ..., 3$ and $l_2 = 0, ..., 3$ in the twovariables case. Observe that the first row and first column of the two-variable case represents the two dimension extension of the one-variable case. In general, Walsh functions are complexvalued, but this is something that we don't have to worry about in base 2.

1. Multiplying two 1-variable elementary Walsh functions is taking the Walsh function on the sum of its digits modulo b (difference, if the second Walsh function is its conjugate). So, for all $k, l \in \mathbb{N}$ and all $x, y \in [0, 1)$, we have

$$\operatorname{wal}_k(x)\operatorname{wal}_l(x) = \operatorname{wal}_{k\oplus_b l}(x),$$

 $\operatorname{wal}_k(x)\overline{\operatorname{wal}_l(x)} = \operatorname{wal}_{k\oplus_b l}(x).$

As the base b is symbolically consistent throughout the analysis, we forgo the subscript to make the arguments look nicer from this point forward.

2. For all $k, l \in \mathbb{N}^s$, the following orthogonality property holds:

$$\int_{[0,1]^s} \operatorname{wal}_{\boldsymbol{k}}(\boldsymbol{x}) \overline{\operatorname{wal}_{\boldsymbol{l}}(\boldsymbol{x})} d\boldsymbol{x} = \begin{cases} 1 & \text{if } \boldsymbol{k} = \boldsymbol{l}, \\ 0 & \text{if } \boldsymbol{k} \neq \boldsymbol{l} \end{cases}$$

3. For any positive integer s, the system $\{\operatorname{wal}_{l}(\boldsymbol{x}) : \boldsymbol{l} \in \mathbb{N}^{s}\}$ is complete and orthonormal in $L^{2}([0,1)^{s})$.

Since the Walsh system is an orthonormal basis for $L^2([0,1)^s)$, each $f \in L^2([0,1)^s)$ has a unique Walsh series decomposition

$$f(\mathbf{x}) \sim \sum_{\boldsymbol{l} \in \mathbb{N}^s} \hat{f}(\boldsymbol{l}) \operatorname{wal}_{\boldsymbol{l}}(\mathbf{x})$$

where \sim denotes the L^2 -equivalence and $\hat{f}(l)$ is the Walsh coefficient of f at l. As a final piece of notation, for each $k \in \mathbb{N}^s$, we let

$$L_{k} := \{ (l_{1}, \dots, l_{s}) \in \mathbb{N}^{s} : \lfloor b^{k_{j}-1} \rfloor \le l_{j} < b^{k_{j}}, \text{ for } j = 1, \dots, s \}$$

and set

$$\sigma_{\boldsymbol{k}}^2(f) = \sum_{\boldsymbol{l} \in L_{\boldsymbol{k}}} |\hat{f}(\boldsymbol{l})|^2.$$

As an example of the usefulness of Walsh coefficients, we conclude this section with a variance result, originally due to Owen [10] using Haar wavelets (see [1] for the Walsh version). More precisely, given an $f \in L^2([0,1)^s)$, the variance of the estimator based on a scrambled (0, m, s)-net can be written as

$$\operatorname{Var}(\hat{I}_n(f)) = \sum_{\boldsymbol{0} \neq \boldsymbol{k} \in \mathbb{N}^s} G_{\boldsymbol{k}} \cdot \sigma_{\boldsymbol{k}}^2(f)$$

where G_k are the gain coefficients. When $G_k < 1$, the scrambled (0, m, s)-net does better than MC for wal_l(\boldsymbol{x}) where $\boldsymbol{l} \in L_k$. When $G_k > 1$, it does worse. There are two key facts about these gain coefficients:

(i)
$$G_{\mathbf{k}} = 0$$
 for $k_1 + \dots + k_s \le m$,
(ii) $G_{\mathbf{k}} \le \left(\frac{b}{b-1}\right)^{\min(s-1,m)} \le e$ for $k_1 + \dots + k_s > m$ (see [10, Theorem 1]).

Using these two properties, one can deduce that the variance of the estimator based on scrambled (0, m, s)-nets converges to 0 faster than MC based on the same number of points. In our work, the Walsh coefficients of the joint pdf take the place of the gain coefficients. When the coefficient $\hat{\psi}(\boldsymbol{l})$ is negative, the RQMC sampling scheme will do better than MC on wal_l(\boldsymbol{x}) and otherwise it will do worse.

3. Walsh Decomposition of the Joint PDF

In this section we write down a formula for the base b Walsh coefficients of the joint pdf of a scrambled (0, m, s)-net in base b. Our calculations will use the fact that the joint pdf is constant on the D_i^s regions. Thus, we will first work towards understanding the Walsh decomposition of $1_{D_i^s}(\mathbf{x}, \mathbf{y})$. Those indicator functions are the product of the two-dimensional functions

$$1_{D_{i_j}}(x_j, y_j) = 1_{C_{i_j}}(x_j, y_j) - 1_{C_{i_j+1}}(x_j, y_j).$$

This brings us to our first lemma.

Lemma 5. The base b Walsh decomposition of the indicator function of C_i is

$$1_{C_i}(x,y) = \sum_{l=0}^{b^i - 1} b^{-i} \operatorname{wal}_l(x) \overline{\operatorname{wal}_l(y)}.$$

Proof. Fix $x, y \in [0, 1)$ and let $\{b^{-i}l\}$ denote the fractional part of $b^{-i}l$. Then

$$\begin{split} \sum_{l=0}^{b^{i}-1} \operatorname{wal}_{l}(x) \overline{\operatorname{wal}_{l}(y)} &= \sum_{l=0}^{b^{i}-1} \operatorname{wal}_{\lfloor b^{i}x \rfloor \ominus \lfloor b^{i}y \rfloor}(\{b^{-i}l\}) \\ &= b^{i} \int_{[0,1)} \operatorname{wal}_{\lfloor b^{i}x \rfloor \ominus \lfloor b^{i}y \rfloor} dt \\ &= \begin{cases} b^{i} & \text{if } \lfloor b^{i}x \rfloor \ominus \lfloor b^{i}y \rfloor = 0, \\ 0 & \text{otherwise.} \end{cases} \end{split}$$

The statement follows because $\lfloor b^i x \rfloor \ominus \lfloor b^i y \rfloor = 0$ exactly when $(x, y) \in C_i$. \Box

The Walsh coefficients of $1_{D_i}(\mathbf{x}, \mathbf{y})$ are found by multiplying together the functions $1_{C_{i_j}}(x_j, y_j) - 1_{C_{i_j+1}}(x_j, y_j)$. In order to keep track of the terms in the product we introduce the following notation.

Definition 6. For $i, l \in \mathbb{N}^s$, denote d(i, l) to be the number of $j \in \{1, \ldots, s\}$ for which $l_j < b^{i_j}$.

Lemma 7. The base b Walsh decomposition of the indicator function of D_i^s is

$$1_{D_i^s}(\mathbf{x}, \mathbf{y}) = \sum_{\boldsymbol{l} \in \mathbb{N}^s} \hat{1}_{D_i^s}(\boldsymbol{l}) \operatorname{wal}_{\boldsymbol{l}}(\mathbf{x}) \overline{\operatorname{wal}_{\boldsymbol{l}}(\mathbf{y})},$$

where

$$\hat{1}_{D_i^s}(l) = \begin{cases} \frac{(-1)^s (1-b)^{d(i,l)}}{b^{s+i}} & \text{if } l_j < b^{i_j+1} \text{ for } j = 1, \dots, s, \\ 0 & \text{otherwise.} \end{cases}$$

Proof. Recall that $D_i = C_i \setminus C_{i+1}$ and $D_i^s = \prod_{j=1}^s D_{i_j}$. This means that

$$\begin{aligned} \mathbf{1}_{D_{i}}(x,y) &= \mathbf{1}_{C_{i}}(x,y) - \mathbf{1}_{C_{i+1}}(x,y) \\ &= \sum_{l=0}^{b^{i}-1} b^{-i} \operatorname{wal}_{l}(x) \overline{\operatorname{wal}_{l}(y)} - \sum_{l=0}^{b^{i+1}-1} b^{-i-1} \operatorname{wal}_{l}(x) \overline{\operatorname{wal}_{l}(y)} \\ &= \sum_{l=0}^{b^{i}-1} \frac{b-1}{b^{i+1}} \operatorname{wal}_{l}(x) \overline{\operatorname{wal}_{l}(y)} - \sum_{l=b^{i}}^{b^{i+1}-1} b^{-i-1} \operatorname{wal}_{l}(x) \overline{\operatorname{wal}_{l}(y)}, \end{aligned}$$

and so

$$\hat{1}_{D_i}(l) = \begin{cases} \frac{(-1)(1-b)^{d(i,l)}}{b^{1+i}} & \text{if } l < b^{i+1}, \\ 0 & \text{otherwise.} \end{cases}$$

Since $1_{D_i^s}(\mathbf{x}, \mathbf{y}) = \prod_{j=1}^s 1_{D_{i_j}}(x_j, y_j)$ and $d(\mathbf{i}, \mathbf{l}) = \sum_{j=1}^s d(i_j, l_j)$, the result follows.

The formula for the Walsh coefficient $\hat{\psi}(l)$ of the joint pdf depends on the non-zero coordinates of l. We call the coordinates on which l is non-zero the support of l and define the following useful vector.

Definition 8. Given $l \in \mathbb{N}^s$ we define its support vector $\operatorname{supp}(l)$ to be the vector r whose j^{th} coordinate is

$$r_j = \begin{cases} 1 & \text{if } l_j > 0, \\ 0 & \text{if } l_j = 0. \end{cases}$$

We can now write the joint pdf using its base b Walsh decomposition.

Proposition 9. Let \tilde{P}_n be a scrambled digital (t, m, s)-net in base b whose projection onto the j^{th} coordinate is a (0, m, s)-net and let $\psi(\mathbf{x}, \mathbf{y})$ be the joint pdf of two distinct points randomly chosen from \tilde{P}_n . Then

(i) The base b Walsh decomposition of $\psi(\mathbf{x}, \mathbf{y})$ takes the form

$$\psi(\mathbf{x}, \mathbf{y}) = 1 + \sum_{\mathbf{l} \in \mathbb{N}^s, \, \mathbf{l} \neq 0} \hat{\psi}(\mathbf{l}) wal_{\mathbf{l}}(\mathbf{x}) \overline{wal_{\mathbf{l}}(\mathbf{y})}, \text{ and }$$

(ii) for $\mathbf{0} \neq \mathbf{l} \in \mathbb{N}^s$, the value of $\hat{\psi}(\mathbf{l})$ in part (i) is

$$\hat{\psi}(l) = \frac{1}{n(n-1)} \left(\frac{b}{b-1}\right)^r \sum_{e \in \{0,1\}^s, \ e \le r} (-1)^e b^{-e} M_b(k-e; \tilde{P}_n),$$
where $r := \operatorname{supp}(l)$ and $k := (|l_1|, \dots, |l_s|).$

Proof. The joint pdf is constant on the D_i^s regions, and we denote these values to be ψ_i . We calculate

$$\begin{split} \hat{\psi}(\boldsymbol{l}) &= \int_{[0,1)^{2s}} \psi(\mathbf{x}, \mathbf{y}) \operatorname{wal}_{\boldsymbol{l}}(\mathbf{x}) \overline{\operatorname{wal}_{\boldsymbol{l}}(\mathbf{y})} d\mathbf{x} d\mathbf{y} \\ &= \int_{[0,1)^{2s}} \sum_{\boldsymbol{i} \in \mathbb{N}^{s}} \psi_{\boldsymbol{i}} \mathbf{1}_{D_{i}^{s}}(\mathbf{x}, \mathbf{y}) \operatorname{wal}_{\boldsymbol{l}}(\mathbf{x}) \overline{\operatorname{wal}_{\boldsymbol{l}}(\mathbf{y})} d\mathbf{x} d\mathbf{y} \\ &= \sum_{\boldsymbol{i} \in \mathbb{N}^{s}} \psi_{\boldsymbol{i}} \hat{\mathbf{1}}_{D_{i}^{s}}(\boldsymbol{l}) \\ &= \sum_{\boldsymbol{i} \in \mathbb{N}^{s}, |l_{j}| - 1 \leq i_{j}} \frac{N_{b}(\boldsymbol{i}; \tilde{P}_{n})}{n(n-1)} \frac{b^{s+i}}{(b-1)^{s}} \frac{(-1)^{s}(1-b)^{d(\boldsymbol{i},\boldsymbol{l})}}{b^{s+i}} \\ &= \frac{1}{n(n-1)} \left(\frac{-1}{b-1}\right)^{s} \sum_{\boldsymbol{i} \in \mathbb{N}^{s}, |l_{j}| - 1 \leq i_{j}} N_{b}(\boldsymbol{i}; \tilde{P}_{n})(1-b)^{d(\boldsymbol{i},\boldsymbol{l})} \\ &= \frac{1}{n(n-1)} \left(\frac{-1}{b-1}\right)^{s} \sum_{\boldsymbol{i} \in \mathbb{N}^{s}, |l_{j}| - 1 \leq i_{j}} \left(\sum_{a=0}^{d(\boldsymbol{i},\boldsymbol{l})} N_{b}(\boldsymbol{i}; \tilde{P}_{n})(-b)^{a} \binom{d(\boldsymbol{i},\boldsymbol{l})}{a}\right). \end{split}$$

Note that the fourth equality holds because $\hat{1}_{D_i^s}(l) = 0$ whenever there is some j for which $|l_j| \ge i_j$. The fifth equality is from Theorem 4.

Next, we will show that

$$\sum_{\boldsymbol{i}\in\mathbb{N}^{s},\,|l_{j}|-1\leq i_{j}} \left(\sum_{a=0}^{d(\boldsymbol{i},\boldsymbol{l})} N_{b}(\boldsymbol{i};\tilde{P}_{n})(-b)^{a} \binom{d(\boldsymbol{i},\boldsymbol{l})}{a} \right)$$
$$= \sum_{\boldsymbol{e}\in\{0,1\}^{s}} \left(\sum_{\boldsymbol{i}\in\mathbb{N}^{s},\,|l_{j}|-e_{j}\leq i_{j}} (-b)^{s-e} N_{b}(\boldsymbol{i};\tilde{P}_{n}) \right).$$

For this to hold it must be that, for a fixed \mathbf{i} and a, the term $(-b)^a N_b(\mathbf{i}; \tilde{P}_n)$ appears exactly $\binom{d(\mathbf{i}, \mathbf{l})}{a}$ times on the right hand side. Therefore we must show that the number of vectors $\mathbf{e} \in \{0, 1\}^s$ such that s - e = a and $|l_j| - e_j \leq i_j$ for $j = 1, \ldots, s$ is $\binom{d(\mathbf{i}, \mathbf{l})}{a}$. To satisfy the second condition, $e_j = 1$ in the $s - d(\mathbf{i}, \mathbf{l})$ coordinates where $|l_j| - 1 = i_j$, leaving $e - s + d(\mathbf{i}, \mathbf{l}) = s - a - s + d(\mathbf{i}, \mathbf{l}) = d(\mathbf{i}, \mathbf{l}) - a$ ones that can be in any of the remaining $d(\mathbf{i}, \mathbf{l})$ coordinates, for which there are indeed $\binom{d(\mathbf{i}, \mathbf{l})}{a}$ possibilities. Using (3), we see that for all $\mathbf{e} \in \{0, 1\}^s$,

$$\sum_{\boldsymbol{i}\in\mathbb{N}^s,\,|l_j|-e_j\leq i_j}(-b)^{s-e}N_b(\boldsymbol{i};\tilde{P}_n)=(-b)^{s-e}M_b(\boldsymbol{k}-\boldsymbol{e};\tilde{P}_n),$$

where $\mathbf{k} = (|l_1|, \dots, |l_s|)$, we may continue our original calculation to obtain

$$\hat{\psi}(l) = \frac{1}{n(n-1)} \left(\frac{b}{b-1}\right)^s \sum_{e \in \{0,1\}^s} (-1)^e b^{-e} M_b(\mathbf{k} - \mathbf{e}; \tilde{P}_n).$$
(4)

Observe that since

$$M_b(\boldsymbol{k}-\boldsymbol{e}; \tilde{P}_n) = M_b(\max(\boldsymbol{k}-\boldsymbol{e}; \boldsymbol{0}); \tilde{P}_n)$$

(the maximum is taken coordinate-wise), the set of values of $M_b(\mathbf{k}-\mathbf{e}; \tilde{P}_n)$ where $\mathbf{e} \in \{0, 1\}^s$ is the same with or without the restriction $\mathbf{e} \leq \mathbf{r} := \operatorname{supp}(\mathbf{k})$. Thus,

$$\sum_{e \in \{0,1\}^s} (-1)^e b^{-e} M_b(\mathbf{k} - \mathbf{e}; \tilde{P}_n)$$

= $\sum_{e \in \{0,1\}^s, e \le r} \left(\sum_{i \in \mathbb{N}^s, i \le 1 - r} (-b)^{-e-i} M_b(\mathbf{k} - \mathbf{e}; \tilde{P}_n) \right)$
= $\sum_{e \in \{0,1\}^s, e \le r} (-b)^{-e} M_b(\mathbf{k} - \mathbf{e}; \tilde{P}_n) \sum_{i=0}^{s-r} (-b)^{-i} {s-r \choose i}$
= $\left(\frac{b-1}{b}\right)^{s-r} \sum_{e \in \{0,1\}^s, e \le r} (-b)^{-e} M_b(\mathbf{k} - \mathbf{e}; \tilde{P}_n).$

Substituting this into (4) completes the proof.

We may simplify the formula for the Walsh coefficients of the joint pdf further in the special case t = 0.

Theorem 10. Let \tilde{P}_n be a scrambled (0, m, s)-net in base b and let $\psi(\mathbf{x}, \mathbf{y})$ be the joint pdf of two distinct points randomly chosen from \tilde{P}_n . Then

(i) the base b Walsh decomposition of $\psi(\mathbf{x}, \mathbf{y})$ takes the form

$$\psi(\mathbf{x}, \mathbf{y}) = 1 + \sum_{l \in \mathbb{N}^s, \ l \neq \mathbf{0}} \hat{\psi}(l) \operatorname{wal}_l(\mathbf{x}) \overline{\operatorname{wal}_l(\mathbf{y})}, \ and$$

(ii) for $\mathbf{0} \neq \mathbf{l} \in \mathbb{N}^s$, the value of $\hat{\psi}(\mathbf{l})$ in part (i) is

$$\hat{\psi}(\mathbf{l}) = -(n-1)^{-1}(1-b)^{1-r} \sum_{i=0}^{r-1-c} (-b)^i \binom{r-1}{i}$$

where r = supp(l) and $c = \max(|l_1| + \dots + |l_s| - m, 0)$.

Proof. Fix $\mathbf{0} \neq \mathbf{l} \in \mathbb{N}^s$ and set $\mathbf{k} = (|l_1|, \dots, |l_s|)$. For a scrambled (0, m, s)-net in base b we have

$$M_b(\mathbf{k} - \mathbf{e}; \tilde{P}_n) = \begin{cases} n(b^{e-c} - 1) & \text{if } e \ge c, \\ 0 & \text{otherwise.} \end{cases}$$

Thus, Proposition 9 (ii) becomes

$$\hat{\psi}(\mathbf{l}) = \frac{1}{n-1} \left(\frac{b}{b-1}\right)^r \sum_{e=c}^r (-1)^e b^{-e} (b^{e-c} - 1) \binom{r}{e}$$
$$= \frac{1}{n-1} \left(\frac{b}{b-1}\right)^r b^{-c} \sum_{j=0}^{r-c} (-1)^{c+j} (1-b^{-j}) \binom{r}{c+j}.$$

From here we make the substitution $1 - b^{-j} = (b - 1) \sum_{i=1}^{j} b^{-i}$ so that

$$\sum_{j=0}^{r-c} (-1)^{c+j} (1-b^{-j}) \binom{r}{c+j} = (b-1) \sum_{j=0}^{r-c} \sum_{i=1}^{j} (-1)^{c+j} b^{-i} \binom{r}{c+j}$$
$$= (b-1) \sum_{i=1}^{r-c} b^{-i} \sum_{j=i}^{r-c} (-1)^{c+j} \binom{r}{c+j},$$

where to change the order of the double sum we observed that i ranges from 1 to r-c and in order for $(-1)^{c+j}b^{-i}\binom{r}{c+j}$ to appear we must have $i \leq j$. Next

we change the index $j \mapsto i + j$ in the inner sum to get

$$\begin{split} \sum_{j=i}^{r-c} (-1)^{c+j} \binom{r}{c+j} &= (-1)^r \sum_{j=0}^{r-c-i} (-1)^{r-c-i-j} \binom{r}{r-c-i-j} \\ &= (-1)^r \sum_{j=0}^{r-c-i} (-1)^j \binom{r}{j} \\ &= (-1)^r (-1)^{r-c-i} \binom{r-1}{r-c-i}, \end{split}$$

where the last equality used a known identity

$$\sum_{i=1}^{k} (-1)^{i} \binom{a}{i} = (-1)^{k} \binom{a-1}{k}.$$

Putting this all together gives

$$\hat{\psi}(\boldsymbol{l}) = -(n-1)^{-1}(1-b)^{1-r} \sum_{i=1}^{r-c} (-b)^{r-c-i} {r-1 \choose r-c-i} = -(n-1)^{-1}(1-b)^{1-r} \sum_{i=0}^{r-1-c} (-b)^{i} {r-1 \choose i}.$$

The previous theorem shows that $\hat{\psi}(\boldsymbol{l})$ depends only on c and the number of non-zero coordinates in \boldsymbol{l} . In the next section it will be helpful to re-index the Walsh coefficients of the joint pdf of a scrambled (0, m, s)-net in base b. Therefore, we make the following definitions.

Definition 11. (Walsh Coefficients)

- 1. Denote $\hat{\psi}_k$ to be the value of the Walsh coefficient $\hat{\psi}(\mathbf{l})$ when $\mathbf{l} \in L_k$.
- 2. For $b, c, r, s \in \mathbb{N}$, $b \ge 2$,

$$\Psi_b^s(r,c) := -(1-b)^{1-r} \cdot \sum_{i=0}^{r-1-c} (-b)^i \binom{r-1}{i}.$$

The second part of the definition gives a covariance equivalent of Owen's gain coefficients [11] that were mentioned in the previous section. However, his analysis only focused on the largest coefficient for which he gave a bound. In the next section, we illustrate that we can do more.

4. Decay Condition on Walsh Coefficients

Using the notation put forth in Section 2.2, Section 2.3 and Theorem 10(i), we obtain the following formula for the covariance term in (2), that is,

$$\operatorname{Cov}(f(\mathbf{U}_I), f(\mathbf{U}_J)) = \sum_{\mathbf{0} \neq \mathbf{l} \in \mathbb{N}^s} |\hat{f}(\mathbf{l})|^2 \hat{\psi}(\mathbf{l}) = \sum_{\mathbf{0} \neq \mathbf{k} \in \mathbb{N}^s} \sigma_{\mathbf{k}}^2(f) \hat{\psi}_{\mathbf{k}}.$$

The remainder of the paper will be devoted to proving that for a particular kind of function, f, this value is less than or equal to 0.

To begin, we must make an assumption on the values of $|\hat{f}(l)|^2$ for $l \in L_k$ or $\sigma_k^2(f)$. Perhaps the most natural conditions we could choose are either

$$|\hat{f}(\boldsymbol{l})|^2 = x^k \alpha_f \quad \text{or} \quad \sigma_{\boldsymbol{k}}^2(f) = x^k \alpha_f$$

s.t. the Walsh series converges (i.e. $x \in [0, b^{-1})$ for the former and $x \in [0, 1)$ for the latter), and α_f is a positive constant that depends on the function. For the purpose of our analysis, we note that α_f can be ignored because multiplication by a positive constant does not change the sign of $\text{Cov}(f(\mathbf{U}_I), f(\mathbf{U}_J))$.

We can rewrite the first decay condition using the fact that $|L_k| = \left(\frac{b-1}{b}\right)^r \cdot b^k$ where r is the number of non-zero coordinates of k as

$$\sigma_k^2(f) = \left(\frac{b-1}{b}\right)^r b^k x^k \alpha_f,\tag{5}$$

which is a less restrictive condition. Both of these decay conditions are a special case of the function $\sigma_k^2(f) = a^r x^k \alpha_f$ which appears in numerous results with an inequality rather than an equality. For example, see [1, Lemma 13.23]. With this formulation, we have evidence to suggest that any values of $a, x \in [0, 1)$ in (5) leads to a covariance term that is not positive, but have been unable to obtain a proof for any case other than for $a = \frac{b-1}{b}$ due to the limitations of the symbolic computation software, as we will show in the last two sections of the paper. Thus, our strategy is to fix a and view the covariance as a polynomial in x of degree m + s - 1 and show that these polynomials are not positive between 0 and 1.

Lemma 12. Let $\tilde{P}_n = {\mathbf{U}_1, \dots, \mathbf{U}_n}$ be a scrambled (0, m, s)-net in base b. Suppose that $f \in L^2([0, 1)^s)$ is a function such that

$$\sigma_{k}^{2}(f) = a^{r}(bx)^{k}\alpha_{f}$$

where α_f is a positive constant that depends on f, $a \in [0,1]$ and $x \in [0,1/b)$ for all $\mathbf{0} \neq \mathbf{k} \in \mathbb{N}^s$ with $k \leq m + r - 1$ where $\mathbf{r} = \operatorname{supp}(\mathbf{k})$. Suppose further that $n = b^m$. Then we can simplify $\frac{b^m - 1}{\alpha_f} \cdot \operatorname{Cov}(f(\mathbf{U}_I), f(\mathbf{U}_J))$ to the polynomial

$$\sum_{k=1}^{m+s-1} \left(\sum_{r=1}^{s} {s \choose r} {k-1 \choose r-1} a^r \Psi_b^s(r, c_m(k)) \right) (bx)^k.$$
(6)

Proof. Define $c_m(k) = \max(k - m, 0)$. We have

$$\begin{split} \frac{b^m - 1}{\alpha_f} \cdot \operatorname{Cov}(f(\mathbf{U}_I), f(\mathbf{U}_J)) &= \frac{b^m - 1}{\alpha_f} \cdot \sum_{\mathbf{0} \neq \mathbf{k} \in \mathbb{N}^s} \sigma_{\mathbf{k}}^2(f) \hat{\psi}_{\mathbf{k}} \\ &= \sum_{\mathbf{0} \neq r \in \{0,1\}^s} \sum_{r \le k \in \mathbb{N}^s} a^r(bx)^k \Psi_m^s(r, c_m(k)) \\ &= \sum_{r=1}^s \binom{s}{r} \sum_{k=r}^{m+r-1} \binom{k-1}{r-1} a^r(bx)^k \Psi_m^s(r, c_m(k)) \\ &= \sum_{k=1}^{m+s-1} \left(\sum_{r=1}^s \binom{s}{r} \binom{k-1}{r-1} a^r \Psi_b^s(r, c_m(k)) \right) (bx)^k \end{split}$$

because there are $\binom{k-1}{r-1}$ ways to partition k into r non-zero parts.

We finish this section with an application of our final result, namely the average case covariance. To begin to do this, we must make sense of integration in $L^2([0,1)^s)$.

Definition 13. Let μ be a Borel probability measure on $L^2([0,1)^s)$ and let

$$F: L^2([0,1)^s) \to L^2([0,1)^s).$$

We say that F is Pettis integrable (weak integrable) with respect to μ if for every continuous linear functional ϕ on $L^2([0,1)^s)$, the complex-valued function

$$\phi \circ F : (L^2([0,1)^s,\mu) \to \mathbb{C}$$

is integrable and there exists some $I(f) \in L^2([0,1)^s)$ such that

$$\phi(I(f)) = \int_{f \in L^2([0,1)^s)} \phi \circ F(f) d\mu(f)$$

holds for all ϕ . In this case, we write

$$I(f) = \int_{f \in L^2([0,1)^s)} F(f) d\mu(f)$$

and say that I(f) is the (Pettis) integral of F with respect to μ .

It is worth a few lines to explain why this definition is useful for us. For a fixed ϕ , we can understand $\int_{f \in L^2([0,1)^s)} \phi \circ F(f) d\mu(f)$ using the standard Lebesgue integral. The Pettis integral simply guarantees that the equation

$$\phi\left(\int_{f\in L^2([0,1)^s)} F(f) \, d\mu(f)\right) = \int_{f\in L^2([0,1)^s)} \phi \circ F(f) \, d\mu(f)$$

makes sense (note that since continuous linear functionals on $L^2([0,1)^s)$ separate points, I(f) must be unique). As an example, consider the function F(f) = f. Then the integral

$$\int_{f\in L^2([0,1)^s)}\phi(f)\,d\mu(f)$$

gives the average value of ϕ with respect to μ . Thus, if there is an $h \in L^2([0,1)^s)$ such that $\phi(h) = \int_{f \in L^2([0,1)^s)} \phi(f) d\mu(f)$, then with respect to continuous linear functionals, h behaves exactly how the mean of μ in $L^2([0,1)^s)$ to behave. In this case, we say that μ admits a mean and call h the mean function of μ .

The following lemma tells us that given a Borel probability measure μ , the average covariance with respect to μ is just the covariance of the mean vector.

Lemma 14. Let μ be a Borel probability measure on $L^2([0,1)^s)$ that admits a mean function h. Let $\tilde{P}_n = \{\mathbf{U}_1, \ldots, \mathbf{U}_n\}$ be a scrambled (t, m, s)-net in base b with joint pdf $\psi(\mathbf{x}, \mathbf{y})$. Then

$$\int_{f \in L^2([0,1)^s)} \operatorname{Cov}(f(\mathbf{U}_I), f(\mathbf{U}_J)) d\mu(f) = \sum_{\mathbf{0} \neq \mathbf{k} \in \mathbb{N}^s} \sigma_{\mathbf{k}}^2(h) \psi_{\mathbf{k}} = \operatorname{Cov}(h(\mathbf{U}_I), h(\mathbf{U}_J))$$

Proof. By definition of the Pettis integral,

$$\begin{split} \int_{f \in L^2([0,1)^s)} \operatorname{Cov}(f(\mathbf{U}_I), f(\mathbf{U}_J)) d\mu(f) &= \int_{f \in L^2([0,1)^s)} \Big(\sum_{l \in \mathbb{N}^s} |\hat{f}(l)|^2 \hat{\psi}(l) \Big) d\mu(f) \\ &= \sum_{l \in \mathbb{N}^s} \left| \Big\langle \int_{f \in L^2([0,1)^s)} f d\mu(f), \operatorname{wal}_l \Big\rangle \Big|^2 \hat{\psi}(l) \\ &= \sum_{l \in \mathbb{N}^s} |\hat{h}(l)|^2 \hat{\psi}(l) \\ &= \sum_{\mathbf{0} \neq \mathbf{k} \in \mathbb{N}^s} \sigma_k^2(h) \psi_k \\ &= \operatorname{Cov}(h(\mathbf{U}_I), h(\mathbf{U}_J)). \end{split}$$

5. Employing Symbolic Computation

We wish to prove that the expected value of the covariance from Lemma 12 is not positive on [0,1) for all $a \in [0,1]$ and $b, m, s \in \mathbb{N} \setminus \{0\}$. For simplicity, we will reduce a parameter, and illustrate how to do this for a = (b-1)/b. This particular choice happens to be well-suited to our computations and is a natural choice of decay to be able to cancel out many of the common factors in the polynomial (6). We remark that a (0, m, s)-net in base b requires $b \ge s - 1$ to exist, whereas the following analysis will not. Figure 4 illustrates the behavior of the polynomials with our chosen a for different values of b, m, s in our domain. To make the pictures nicer, we include the scaling factor $(b^m - 1)^{-1}$, which doesn't modify the sign of the polynomial and still gives an impression of a general pattern. In particular, 4(c) is only an observation for the polynomials themselves and will not make sense for our main results.

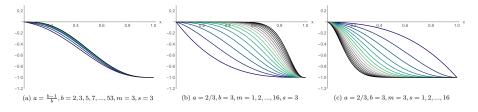


Figure 4: A scaled version of polynomial (6) for different values of b, m, s and $a = \frac{b-1}{b}$ where shades of blue represent smaller values of the varying parameter (darkest is smallest) and shades of gray represent larger values (darkest is largest).

Creative Telescoping. We employ the principle of creative telescoping [15] in an attempt to further simplify polynomial (6) effectively. To see how this could be possible, we first substitute the formula for $\Psi_b^s(r,c)$ from Definition 11 into (6) and fix a = (b-1)/b to get

$$\sum_{k=1}^{m+s-1} \left(\sum_{r=1}^{s} \binom{s}{r} \binom{k-1}{r-1} \frac{b-1}{(-b)^r} \sum_{i=0}^{r-1-c_m(k)} (-b)^i \binom{r-1}{i} \right) (bx)^k.$$
(7)

We first note that the summands of this triple sum contain holonomic functions in the parameters (roughly speaking, the binomial coefficients and exponential functions satisfy recurrences with polynomial coefficients), whose products and sums are also holonomic [4, Theorem 2.16]. As an additional simplification, the innermost sum can be split into two cases based on $c_m(k)$, both of which collapse into double sums with summands that are holonomic. We can then invoke the function CREATIVETELESCOPING from the package HOLONOMICFUNCTIONS.M [5] which provides telescoping relations for our multiple sums. With some careful manipulation to treat issues of singularities and unnatural boundary values, along with a tedious ~ 30 hours of computation, we managed to obtain a recurrence that (7) satisfies. We encourage the reader to refer to the accompanying Mathematica notebook for these computations, which can be downloaded from website mentioned in the introduction. In this way, we are able to assert the following lemma. **Lemma 15.** For $b, m, s \in \mathbb{N}, b \geq 2$, the polynomial (6) satisfies the recurrence

$$\begin{split} &(s+2)(bx-1)\cdot c[s+3]\\ &+(m(bx-1)(x-1)+bsx(x-2)+bx(x-3)-s(2x-3)-3x+5)\cdot c[s+2]\\ &-(x-1)(bmx+bsx+bx+mx-2m+sx-3s+x-4)\cdot c[s+1]\\ &+(x-1)^2(m+s+1)\cdot c[s]\\ &=0. \end{split}$$

This recurrence from the above lemma can be solved using the SOLVERECURRENCE command from the SIGMA.M [12] package. We were able to obtain a non-trivial solution with the following form as an output:

$$(1 - (bx)^{m}) - (1 - (bx)^{m}) \left(\frac{x-1}{bx-1}\right)^{s} + (bx)^{m} \left(\frac{x^{m}-1}{x^{m}} + \frac{\Gamma(m+s+1)}{\Gamma(m)\Gamma(s+2)}(1-x)^{s+1}{}_{2}F_{1}\left[\frac{m+s+1,1}{s+2};1-x\right]\right)$$
(8)
+ $(bx)^{m} \left(\frac{x-1}{bx-1}\right)^{s} \left(\frac{1-(bx)^{m}}{(bx)^{m}} - \frac{\Gamma(m+s+1)}{\Gamma(m)\Gamma(s+2)}(1-bx)^{s+1}{}_{2}F_{1}\left[\frac{m+s+1,1}{s+2};1-bx\right]\right)$

The Digital Library of Mathematical Functions [2] provides a list of identities that allows us to make nice simplifications, such as using the Beta function in DLMF 8.17.8:

$$_{2}F_{1}\begin{bmatrix}a+b,1\\a+1\end{bmatrix} = \frac{a}{x^{a}(1-x)^{b}}B_{x}(a,b)$$

where

$$B_x(a,b) := \int_0^x t^{a-1} (1-t)^{b-1} dt.$$

In particular,

$$B_1(a,b) = B(a,b) = \frac{\Gamma(a)\Gamma(b)}{\Gamma(a+b)}.$$

Using this with a = s + 1 and b = m and simplifying the fractions containing $\Gamma(n) = (n - 1)!$, our polynomial now becomes:

$$(1 - (bx)^m) - (1 - (bx)^m) \left(\frac{x - 1}{bx - 1}\right)^s + (bx)^m \left(\frac{x^m - 1}{x^m} + \frac{1}{x^m} \frac{B_{1-x}(s + 1, m)}{B(s + 1, m)}\right) + (bx)^m \left(\frac{x - 1}{bx - 1}\right)^s \left(\frac{1 - (bx)^m}{(bx)^m} - \frac{1}{(bx)^m} \frac{B_{1-bx}(s + 1, m)}{B(s + 1, m)}\right).$$

In the expression above, we observe the presence of the normalized beta function $I_x(a,b) = B_x(a,b)/B(a,b)$. By DLMF 8.17.4, we have the identity

$$I_x(a,b) = 1 - I_{1-x}(b,a).$$

This now gives:

$$\begin{split} &(1-(bx)^m) - (1-(bx)^m) \left(\frac{x-1}{bx-1}\right)^s \\ &+ (bx)^m \left(\frac{x^m-1}{x^m} + \frac{1}{x^m}(1-I_x(m,s+1))\right) \\ &+ (bx)^m \left(\frac{x-1}{bx-1}\right)^s \left(\frac{1-(bx)^m}{(bx)^m} - \frac{1}{(bx)^m}I_{1-bx}(s+1,m)\right). \end{split}$$

A final simplification gives:

$$Q_s(b,m,x) := 1 - b^m I_x(m,s+1) - \left(\frac{1-x}{1-bx}\right)^s I_{1-bx}(s+1,m).$$

Lemma 15 reveals that $Q_s(b, m, x)$ is a sign equivalent formulation for (7), so we use it for the remainder of our analysis. The main result of this section (i.e. the fact that $Q_s(b, m, x)$ is not positive on [0, 1)) is derived using properties found in Section 8 of the DLMF [2]. However, one interesting property that we need cannot be found there, so we conjecture and prove it in the lemma below.

Lemma 16. The normalized Beta function can be simplified to a derivative function as follows

$$I_x(a,b) = \frac{(-x)^a}{(a-1)!} \cdot D_x^{a-1}\left(\frac{(1-x)^{a+b-1}-1}{x}\right),$$

where $D_x^j(\cdot)$ is the *j*-th partial derivative of the expression (\cdot) with respect to the variable x. The formula holds for $x \in [0, 1], a, b \in \mathbb{N}$.

Proof.

$$\begin{split} I_x(a,b) &= \frac{1}{B(a,b)} \cdot \int_0^x t^{a-1} \cdot (1-t)^{b-1} dt \\ &= \frac{(a+b-1)!}{(a-1)! \cdot (b-1)!} \cdot \int_0^x \sum_{i=0}^{b-1} \binom{b-1}{i} \cdot (-1)^i \cdot t^{i+a-1} dt \\ &= \frac{1}{(a-1)!} \cdot \sum_{i=0}^{b-1} \frac{(a+b-1)!}{(b-1-i)! \cdot i!} \cdot (-1)^i \cdot \frac{x^{i+a}}{i+a} \\ &= \frac{x^a}{(a-1)!} \cdot \sum_{i=a}^{a+b-1} \binom{a+b-1}{i} \cdot \frac{(i-1)!}{(i-a)!} \cdot (-x)^{i-a}, \end{split}$$

which we can write as a derivative

$$I_x(a,b) = \frac{(-1)^{a-1} \cdot x^a}{(a-1)!} \cdot D_x^{a-1} \left(\sum_{i=1}^{a+b-1} \binom{a+b-1}{i} \cdot (-x)^{i-1} \right)$$
$$= -\frac{(-x)^a}{(a-1)!} \cdot D_x^{a-1} \left(\frac{(1-x)^{a+b-1}-1}{-x} \right)$$
$$= \frac{(-x)^a}{(a-1)!} \cdot D_x^{a-1} \left(\frac{(1-x)^{a+b-1}-1}{x} \right).$$

We now proceed to show the main result of this section.

Theorem 17. For $b, m, s \in \mathbb{N}, b \ge 2, 0 \le x \le 1$,

$$Q_s(b, m, x) \le 0,$$

where

$$Q_s(b,m,x) := 1 - b^m I_x(m,s+1) - \left(\frac{1-x}{1-bx}\right)^s I_{1-bx}(s+1,m).$$

Proof. We use an inductive style proof on s. First, we verify for s = 1 that

$$Q_1(b,m,x) = (1-b)x \cdot \frac{(bx)^m - 1}{bx - 1} = (1-b)x \cdot \sum_{i=0}^{m-1} (bx)^i \le 0$$

Induction tells us that it is enough to show that $Q_{s-1} \leq 0 \Rightarrow Q_s \leq 0$, but instead, we choose to show

$$\Delta_s := Q_{s-1} - Q_s \ge 0,$$

which would imply that the polynomials are decreasing as a function of s, thereby giving us our result. We first observe that Δ_s can be separated into two parts, $\Delta_s^{(1)}$ and $\Delta_s^{(2)}$ (because the 1's cancel in the difference), and we simplify each of the two parts separately.

 $\underline{Part 1:}$

$$\Delta_s^{(1)} := b^m (I_x(m, s+1) - I_x(m, s))$$

To simplify this, we look at the integral representation of the beta function.

$$\begin{split} \Delta_s^{(1)} &= \frac{b^m}{B(m,s+1)} \cdot \int_0^x t^{m-1} (1-t)^s dt - \frac{b^m}{B(m,s)} \cdot \int_0^x t^{m-1} (1-t)^{s-1} dt \\ &= b^m \cdot \frac{(m+s-1)!}{(m-1)! \cdot (s-1)!} \cdot \int_0^x \left(\frac{m+s}{s} \cdot t^{m-1} \cdot (1-t)^s - t^{m-1} \cdot (1-t)^{s-1} \right) dt \\ &= b^m \cdot \frac{(m+s-1)!}{(m-1)! \cdot s!} \cdot \int_0^x \left(mt^{m-1} \cdot (1-t)^s - t^m \cdot s(1-t)^{s-1} \right) dt \\ &= b^m \cdot \binom{m+s-1}{s} \cdot \int_0^x \left(t^m \cdot (1-t)^s \right)' dt \\ &= b^m \cdot \binom{m+s-1}{s} \cdot x^m \cdot (1-x)^s \end{split}$$

<u>Part 2:</u> We write $\Delta_s^{(2)}$ as the difference $\delta_s - \delta_{s-1}$ where

$$\delta_s := (1-x)^s \cdot \frac{I_{1-bx}(s+1,m)}{(1-bx)^s}.$$

For this simplification, we take advantage of Lemma 16. Upon substitution, the derivative simplifies nicely to a symbolic sum in terms of only bx (and not 1-bx for example). We remark that in this case, the use of the identity DLMF 8.17.4 is less elegant. Instead, we simplify as follows:

$$\begin{split} \delta_s &= \frac{(1-x)^s \cdot (-(1-bx))^{s+1}}{s! \cdot (1-bx)^s \cdot (-b)^s} \cdot D_x^s \left(\frac{(1-(1-bx))^{s+m}-1}{1-bx}\right) \\ &= \frac{(1-x)^s \cdot (1-bx)}{s! \cdot b^s} \cdot D_x^s \left(\frac{(bx)^{s+m}-1}{bx-1}\right) \\ &= \frac{(1-x)^s \cdot (1-bx)}{s! \cdot b^s} \cdot D_x^s \left(\sum_{i=0}^{s+m-1} (bx)^i\right) \\ &= \frac{(1-x)^s \cdot (1-bx)}{s! \cdot b^s} \cdot \sum_{i=s}^{s+m-1} \frac{i!}{(i-s)!} \cdot b^i \cdot x^{i-s} \\ &= \frac{(1-x)^s \cdot (1-bx)}{s! \cdot b^s} \cdot \sum_{i=0}^{m-1} \frac{(i+s)!}{i!} \cdot b^{i+s} \cdot x^i \\ &= (1-x)^s \cdot (1-bx) \cdot \sum_{i=0}^{m-1} \binom{i+s}{s} \cdot (bx)^i. \end{split}$$

Combining, $\Delta_s = \Delta_s^{(1)} + \Delta_s^{(2)} = \Delta_s^{(1)} + \delta_s - \delta_{s-1} = (1-x)^s \cdot d_s - \delta_{s-1}$ with

$$\begin{split} d_{s} &= \binom{m+s-1}{s} \cdot (bx)^{m} + (1-bx) \cdot \sum_{i=0}^{m-1} \binom{i+s}{s} \cdot (bx)^{i} \\ &= \binom{m+s-1}{s} \cdot (bx)^{m} + \sum_{i=0}^{m-1} \binom{i+s}{s} \cdot ((bx)^{i} - (bx)^{i+1}) \\ &= \binom{m+s-1}{s} \cdot (bx)^{m} - \sum_{i=1}^{m} \binom{i+s-1}{s} \cdot (bx)^{i} + \sum_{i=0}^{m-1} \binom{i+s}{s} \cdot (bx)^{i} \\ &= \sum_{i=0}^{m-1} \binom{i+s-1}{s-1} \cdot (bx)^{i}, \end{split}$$

gives us the nice formula

$$\Delta_s = (1-x)^s \cdot d_s - \delta_{s-1} = x \cdot (b-1) \cdot (1-x)^s \cdot \sum_{i=0}^{m-1} \binom{i+s-1}{s-1} \cdot (bx)^i,$$

and this is clearly positive for our assumed values, as desired.

Finally, we can conclude with the main result of the paper.

Theorem 18. Let \tilde{P}_n be a scrambled (0, m, s)-net in base b and let $f \in L^2([0, 1)^s)$ be a function whose base b Walsh series decomposition satisfies

$$\sigma_k^2(f) = \left(\frac{b-1}{b}\right)^r (bx)^k \alpha_f,$$

where α_f is a positive constant that depends on $f, x \in [0, 1/b), k = k_1 + \cdots + k_s$ and r is the number of non-zero coordinates of $\mathbf{k} \in \mathbb{N}^s$. Then

$$\operatorname{Cov}(f(\mathbf{U}_I), f(\mathbf{U}_J)) \leq 0 \text{ and } \operatorname{Var}(\hat{I}_n(f)) \leq \operatorname{Var}(\hat{I}_{MC,n}(f)).$$

Proof. From Lemma 15, we deduce that $Q_s(b, m, x)$ gives us the polynomial that is sign equivalent to (6) in Lemma 12 for all parameters in the desired ranges. By Theorem 17, $Q_s(b, m, x) \leq 0$. Thus, we conclude that $\text{Cov}(f(\mathbf{U}_I), f(\mathbf{U}_J)) \leq 0$. Then (1) implies $\text{Var}(\hat{I}_n(f)) \leq \text{Var}(\hat{I}_{MC,n}(f))$.

Corollary 19. Let \tilde{P}_n be a scrambled (0, m, s)-net in base b and suppose μ is a positive Borel probability measure on $L^2([0, 1)^s)$ that admits a mean function h. If the base b Walsh series decomposition of h satisfies

$$\sigma_{\boldsymbol{k}}^2(h) = \left(\frac{b-1}{b}\right)^r (bx)^k \alpha_h,$$

where is α_h positive constant that depends on $h, x \in [0, 1/b), k = k_1 + \cdots + k_s$ and r is the number of non-zero coordinates of $\mathbf{k} \in \mathbb{N}^s$. Then

$$\int_{f \in L^2([0,1)^s)} \operatorname{Cov}(f(\mathbf{U}_I), f(\mathbf{U}_J)) d\mu(f) \le 0.$$

Proof. Applying Lemma 14 and the previous theorem gives

$$\int_{f \in L^2([0,1]^s)} \operatorname{Cov}(f(\mathbf{U}_I), f(\mathbf{U}_J)) d\mu(f) = \operatorname{Cov}(h(\mathbf{U}_I), h(\mathbf{U}_J)) \le 0. \qquad \Box$$

6. Conclusions and Future Work

In this paper we computed the base b Walsh series decomposition of the joint pdf of a scrambled (0, m, s)-net in base b. This allowed us to give a formula for $\operatorname{Cov}(f(\mathbf{U}_I), f(\mathbf{U}_J))$ in terms of the function's Walsh coefficients. Using symbolic computation we were able to show that with a reasonable assumption on the base b Walsh coefficients of f, the covariance term will be negative. This work extends the list of functions for which we know that an estimator based on a scrambled (0, m, s)-net in base b will do no worse than Monte Carlo. This paper also shows that symbolic computation can be applied to QMC integration problems. In future work, we would like to find functions that satisfy the decay condition of Theorem 18, or find a practical measure on $L^2([0, 1)^s)$ whose mean function satisfies the same condition.

Figures 5 and 6 give some insight as to why a more generalized result could not be so easily proved. The nice patterns that we had observed in Figure 4 with $a = \frac{b-1}{b}$ are not as regular here, with polynomials with a closer to 1 exhibiting more erratic behavior in the interval [0, 1).

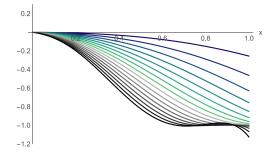


Figure 5: A scaled version of polynomial (6) for b = 3, m = 3, s = 3, a = 1/16, ..., 1 with shades of blue representing smaller values of a (darkest is smallest) and shades of gray representing larger values of a (darkest is largest).

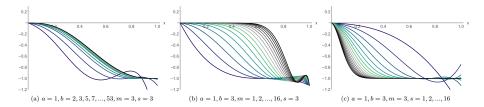


Figure 6: A scaled version of polynomial (6) for a = 1 and different values of b, m, s, where shades of blue represent smaller values of the varying parameter (darkest is smallest) and shades of gray represent larger values (darkest is largest).

We conclude with a conjecture that generalizes our result from Theorem 18 based on experimental evidence. Unfortunately, our guess and then prove technique turned out to be ineffective for all other a except for (b-1)/b, and we were unable to find non-trivial solutions for the recurrences that we were able to obtain for such a. However, as shown in Figures 5 and 6, there is reasonable evidence to show that the result holds.

Conjecture 20. Let \tilde{P}_n be a scrambled (0, m, s)-net in base b and $f \in L^2([0, 1)^s)$ be a function whose base b Walsh series decomposition satisfies

$$\sigma_k^2(f) = a^r x^k \alpha_f,$$

where $a \in [0,1]$, $x \in [0,1)$, $k = k_1 + \cdots + k_s$, α_f is a positive constant that depends on f and r is the number of non-zero coordinates of $\mathbf{k} \in \mathbb{N}^s$. Then

$$\operatorname{Cov}(f(\mathbf{U}_I), f(\mathbf{U}_J)) \leq 0 \text{ and } \operatorname{Var}(I_n(f)) \leq \operatorname{Var}(I_{MC,n}(f)).$$

Acknowledgements. We are particularly grateful to Josef Dick, Christoph Koutschan, Peter Kritzer and Christiane Lemieux for taking time out of their busy schedules to guide us in the right direction at the beginning, and their subsequent encouragement towards the completion of this work. Both authors want to especially acknowledge Christoph for his valuable comments that improved this manuscript greatly. E. Wong would also like to thank Manuel Kauers and Veronika Pillwein for the opportunity to give a talk about this work at OPSFA and to Lin Jiu, Mehdi Makhul, Isabel Pirsic and Ali Uncu for some helpful commentary. E. Wong is supported by the Austrian Science Fund (FWF): F5011-N15. J. Wiart is supported by the Austrian Science Fund (FWF), Projects F5506-N26 and F5509-N26, which are parts of the Special Research Program "Quasi-Monte Carlo Methods: Theory and Applications".

References

- [1] J. Dick, F. Pillichshammer, Digital nets and sequences: Discrepancy theory and quasi-monte carlo integration, Cambridge University Press, UK, 2010.
- [2] DLMF, NIST digital library of mathematical functions, 2019. F. W. J. Olver, A. B. Olde Daalhuis, D. W. Lozier, B. I. Schneider, R. F. Boisvert, C. W. Clark, B. R. Miller and B. V. Saunders, eds., http://dlmf.nist.gov/.
- [3] M. Kauers, Guessing handbook, Technical Report 09-07, RISC Report Series, Johannes Kepler University, Linz, Austria, 2009. Http://www.risc.jku.at/research/combinat/software/Guess/.
- [4] C. Koutschan, Advanced applications of the holonomic systems approach, Ph.D. thesis, Johannes Kepler University, Linz, Austria, 2009.
- [5] C. Koutschan, HolonomicFunctions user's guide, Technical Report 10-01, RISC Report Series, Johannes Kepler University, Linz, Austria, 2010. Http://www.risc.jku.at/publications/download/risc_3934/hf.pdf.
- [6] C. Lemieux, Negative dependence, scrambled nets, and variance bounds, Mathematics of Operations Research 43 (2017) 228–251.
- [7] H. Niederreiter, Low-discrepancy point sets obtained by digital constructions over finite fields, Czechoslovak Math. J 42 (1992) 143–166.
- [8] H. Niederreiter, Random number generation and quasi-Monte Carlo methods, volume 63, SIAM CBMS-NSF Regional Conference Series in Applied Mathematics, 1992.
- [9] A.B. Owen, Randomly permuted (t, m, s)-nets and (t, s)-sequences, in: H. Niederreiter, P.J. Shiue (Eds.), Monte Carlo and Quasi-Monte Carlo Methods in Scientific Computing, volume 106, Springer, New York, NY, 1995, pp. 299–317.
- [10] A.B. Owen, Scrambled net variance for integrals of smooth functions, The Annals of Statistics 25 (1997) 1541–1562.
- [11] A.B. Owen, Variance and discrepancy with alternative scramblings, ACM Transactions on Modeling and Computer Simulation 13 (2003) 363–378.
- [12] C. Schneider, Symbolic summation assists combinatorics, Séminaire Lotharingien de Combinatoire 56 (2007) 1–36. Article B56b, http://www.risc.jku.at/research/combinat/software/Sigma/.
- [13] J.L. Walsh, A closed set of normal orthogonal functions, Amer. J. Math. 45 (1922).
- [14] J. Wiart, C. Lemieux, G. Dong, On the dependence structure of scrambled (t, m, s)-nets, arXiv e-prints (2019).
- [15] D. Zeilberger, The method of creative telescoping, Journal of Symbolic Computation 11 (1991) 195–204.