

Localization and the Noncommutative Mikusiński Calculus

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Joint Work with A. Korpöral (RISC)
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Applications of Computer Algebra (ACA 2012)
25–28 June, 2012, Sofia, Bulgaria



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Definition (Rosenkranz 2003/ Rosenkranz-Regensburger 2006)

We call $(\mathcal{F}, \partial, \int)$ an **integro-differential algebra** over K if \mathcal{F} is a commutative K -algebra with K -linear operations ∂ and \int satisfying

$$(\int f)' = f,$$

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For every $(\mathcal{F}, \partial, \int)$ with characters $\Phi = \{L, R, \dots\}$, there is a computable ring of **integro-differential operators** $\mathcal{F}[\partial, \int] = \mathcal{F}[\partial] \dot{+} \mathcal{F}[\int] \dot{+} (\Phi)$.

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In particular, $D \in \mathcal{F}[\partial]$ and $A, B \in \mathcal{F}[\int]$ for standard case $\mathcal{F} = C^\infty(\mathbb{R})$.

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Theorem (Factorization Theorem)

Given $(T, \mathcal{B}) \in \mathcal{F}_0[\partial]_\Phi$, any factorization $T = T_1 T_2$ lifts to a problem factorization $(T, \mathcal{B}) = (T_1, \mathcal{B}_1) \cdot (T_2, \mathcal{B}_2)$ with $(T_1, \mathcal{B}_1), (T_2, \mathcal{B}_2) \in \mathcal{F}_0[\partial]_\Phi$ and $\mathcal{B}_2 \leq \mathcal{B}$.

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Lemma (Division Lemma)

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Let Φ be an umbral character set for \mathcal{F} . Then for any $(T, \mathcal{B}) \in \mathcal{F}_0[\partial]_{(\Phi)}$ there is a $(S, \mathcal{A}) \in \mathcal{F}_0[\partial]_\Phi$ that has (T, \mathcal{B}) as a subproblem.

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Theorem

Let M be a left R -module, and let $S \subseteq R$ be a multiplicative, right permutable and right reversible denominator set $S \subseteq R$. Then there exists a left $S^{-1}R$ -module $S^{-1}M$.

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Let $(\mathcal{F}, \partial, \int)$ be an integro-differential algebra with character set Φ .
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THANK YOU



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