

Approximate Commutative Algebra – an Impossible Concept ?

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Algebra and **Approximation** a Fundamental Antithesis

Algebra

discrete world

symbols, structures

no neighborhoods

equal $\langle \Rightarrow \rangle$

not equal

satisfied $\langle \Rightarrow \rangle$

not satisfied

true $\langle \Rightarrow \rangle$ false

discrete

elements

equations

relations

mappings

Analysis

continuous world

spaces, functions

neighborhoods

equal ... close

... distant

sat. ... nearly sat.

... not satisfied

nearly true ... false

continuous

Approximation can only be considered

in an **analytic** context, hence

(Commutative) Algebra and **Approximation**

are **antithetic** concepts

But many problems in Commutative Algebra are posed over number fields which carry a

natural topology

e.g. real numbers or complex numbers

Here, $1.41421 \approx \sqrt{2}$ may be
an obvious **approximation**

This suggests an

embedding

of **commutative algebra** into **analysis**
via the coefficient field

Successful Prototype for this Embedding :

Linear Algebra over \mathbb{R}^n or \mathbb{C}^n

becomes

Analytic Geometry

becomes

Numerical Linear Algebra

Embedding into **Analysis** permits

- introduction of **approximation**
- handling of data with **limited accuracy**
- use of **analytic tools** for problem solving
- use of **floating-point** in computation

Embedding requires

- **reconsideration** of all concepts
- **redefinition** of many concepts
- introduction of various **new concepts**
- **disposal** of some classical concepts

Embedding provides

- new **insights**
- means to **avoid ill-conditioning**
- **faster** algorithms
- **safer** algorithms

These benefits arise even for **exact** computation with **exact** data !

Analytic treatment of an **algebraic** problem
has been introduced by

C. F. Gauss

Solution of systems of linear equations
by **iterative approximation**

A Revolutionary Idea :

Have an **algebraic** problem
whose (exact) solution may be obtained by
a *finite number* of (exact) *rational operations*

Instead take an approximate solution and
improve it iteratively and

Stop when **accuracy is sufficient**
e.g. relative to *data accuracy*

Developed and used by Gauss
in a large-scale surveying project
Communicated (by letter) to another surveyor
on Dec. 26, 1823

Numerous details of the **numerical** computation
are touched upon in this communication :

- computation of a good initial approximation
- work with *corrections* not with values
- meaningful accuracy during iteration
- details of choosing the iteration sequence
- stopping criterion

Today , in Scientific Computing ,

- “all” linear systems solved in **floating-point**
- “almost all” linear systems solved **iteratively**

where “solved” \equiv solved **approximately**

Is this **meaningful and **achievable**
also for **polynomial problems** ?**

Fundamental Idea for the **Analytic Treatment** of Quantitative **Algebraic** Problems

“quantitative”: Besides by their structure, problem and results are characterized by **numbers**

Problem : “data” $a \in \mathcal{A} \subseteq \mathbb{C}^M$

Result : “solution” $z \in \mathcal{Z} \subseteq \mathbb{C}^m$

Model of Quantitative Problem :

Given: a **mapping** $F : \mathcal{A} \rightarrow \mathcal{Z}$ (implicitly!)

Sought: $F(a)$ for specified $a \in \text{dom}F \subset \mathcal{A}$

F is the **exact** “data→result mapping”

structural aspects of problem are contained in F

definition of F may include values like 1, 0, etc.

must include structural sparsity

Note: a coefficient 0 may signify

- an absent term (*sparsity*)

- a tiny value (*approximate data*)

Analytic Properties of $F : \mathcal{A} \rightarrow \mathcal{Z}$

$$\dim \mathcal{A} =: M, \quad \dim \mathcal{Z} =: m$$

F is **well-posed** in a neighborhood $N(a)$ of a if

domain $F \supset N(a)$, *open* in \mathcal{A}

$m \leq M$ and F is *surjective* in $N(a)$

F is *continuous* in $N(a)$

otherwise it is **ill-posed**

F is **well-conditioned** in $N(a)$ if

F is *well-posed* in $N(a)$

F is *differentiable* in $N(a)$

$\|F'\|$ is “small”

F is **ill-conditioned** if $\|F'\|$ “large” in $N(a)$

Ill-conditioning implies that the solution z is *highly sensitive* to certain changes in a

Ill-posed Algebraic Problems

Type 1 : $\dim (\text{domain of } F) =: d < M$

Results only defined for *data on a manifold* \mathcal{S}
of dimension d in \mathcal{A} (i.e. only for special data)

Type 2 : $\dim (\text{image of } F) =: \bar{d} < m$

All exact *results* lie on a *manifold* $\bar{\mathcal{S}}$
of dimension \bar{d} in \mathcal{Z}

(i.e. **approximate** solution is **not** a proper result)

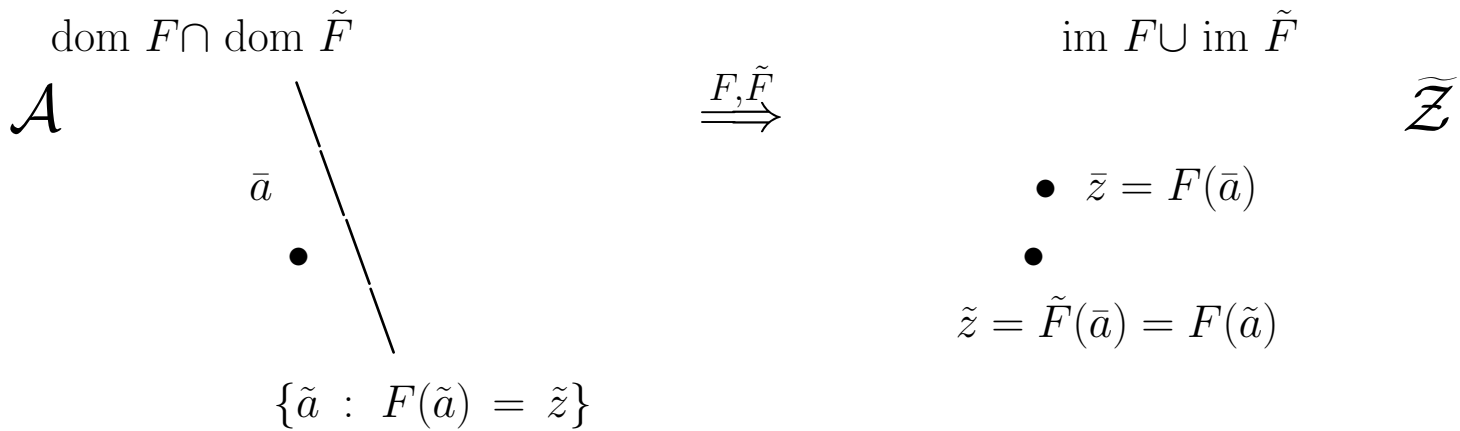
Type 3 : F is *discontinuous* in $N(a)$

i.e. results “jump” at nearby data

Many classes of interesting quantitative
algebraic problems are **ill-posed**

Meaningful *approaches* to **ill-posed** problems are
an important task of Approximate Algebra

Approximate Commutative Algebra



Tasks :

- design **approximate map** \tilde{F}
(algorithm, implementation)
- evaluate \tilde{F} for $\bar{a} \Rightarrow \tilde{z}$
(run code)
- find \tilde{a} **closest to** \bar{a} , with $F(\tilde{a}) = \tilde{z}$
(**backward error**)
- correct \tilde{z} towards \bar{z}
(**successive approximation**)
- find **analytic** properties of F around \bar{a}
(**wellposedness, condition, nearsingularity**)
- re-define F and/or \tilde{F} if appropriate
(respect **analytic** properties of F)

Topologies in \mathcal{A} and \mathcal{Z}

Formally, we have

$$\text{data space } \mathcal{A} \equiv \mathbb{C}^M$$

$$\text{result space } \mathcal{Z} \equiv \mathbb{C}^m$$

Note: **Structural** data are **not** components of \mathcal{A}

May use some **weighted** norm to define distance:

$$\text{2-norm : } \quad \| \tilde{a} - a \|_{e,2} := \left(\sum_{\nu} \left(\frac{|\tilde{\alpha}_{\nu} - \alpha_{\nu}|}{\varepsilon_{\nu}} \right)^2 \right)^{1/2}$$

$$\text{max-norm : } \quad \| \tilde{a} - a \|_{e,\infty} := \max_{\nu} \frac{|\tilde{\alpha}_{\nu} - \alpha_{\nu}|}{\varepsilon_{\nu}}$$

Weights ε_{ν} should relate to **accuracy levels** of α_{ν}

e.g. $\bar{\alpha}_{\nu}$ specified to 4 decimals $\Rightarrow \varepsilon_{\nu} = 10^{-4}$

$$e := (\varepsilon_1, \dots, \varepsilon_M)$$

Different data components α_{ν}

may have different **tolerances** ε_{ν}

Empirical Data

data from measurements, observations etc. have

limited accuracy

specified data $\bar{a} \in \mathcal{A} \subset \mathbb{R}^M$ or \mathbb{C}^M

tolerance $e \in \mathbb{R}_+^M$

validity parameter $\delta \in \mathbb{R}_+$

family of sets of **equivalent data** ($\delta > 0$) :

$$N_\delta(\bar{a}, e) := \{\tilde{a} \in \mathcal{A} : \|\tilde{a} - \bar{a}\|_e \leq \delta\}$$

validity scale (*continuous gradual transition!*)

δ : 0 ... 1 ... 3 ... 10 ... 30 ...
valid probably possibly probably invalid
 valid valid invalid

Empirical polynomial in \mathbb{C}^s :

$$p(x; a) := \sum_{\nu} \alpha_{\nu} x^{j_{\nu}} \quad \text{with} \quad x^{j_{\nu}} := x_1^{j_{\nu 1}} \dots x_s^{j_{\nu s}}$$

where *some* (or all) coefficients α_{ν} are **empirical**

Valid Instances of Empirical Data

$\tilde{a} \in \mathcal{A}$ is a **valid instance** for (\bar{a}, e) if

$$\tilde{a} \in N_\delta(\bar{a}, e) \quad \text{with } \delta = \text{“O(1)”}$$

i.e. \tilde{a} is *indistinguishable* from \bar{a}
on the specified accuracy level,

Valid Results for Empirical Data

computational result $\tilde{z} = \tilde{F}(\bar{a}) \in \mathcal{Z}$

interpreted as **exact** result for **modified** data \tilde{a}

Equivalent-data manifold $\mathcal{M}(\tilde{z})$:

$$\mathcal{M}(\tilde{z}) := \{ \tilde{a} \in \mathcal{A} : \tilde{a} \xrightarrow{F} \tilde{z} \} \subset \mathcal{A}$$

E.g. Polynomial zeros :

$$p(x; \bar{a}) := \bar{\alpha}_0 + \bar{\alpha}_1 x + \dots + \bar{\alpha}_{M-1} x^{M-1} + x^M$$

$$\mathcal{M}(\tilde{z}) := \{ \tilde{a} \in \mathcal{A} : p(\tilde{z}; \tilde{a}) = 0 \} \quad \text{linear (M-1)-manifold}$$

$$\tilde{z} \text{ } \delta\text{-valid} \iff N_\delta(\bar{a}, e) \cap \mathcal{M}(\tilde{z}) \neq \emptyset$$

Find *minimal* δ : \tilde{z} is **valid** if $\delta_{\min} = \text{O(1)}$

Validity Checking

Have to **compute**

$$\delta_{\min}(\tilde{z}) := \min (\|\tilde{a} - \bar{a}\|_e, \tilde{a} \in \mathcal{M}(\tilde{z}))$$

For a *linear* manifold $\mathcal{M}(\tilde{z})$, this is a

- linear *least squares* problem with a *2-norm*
- linear *optimization* problem with a *max-norm*

E.g. Polynomial zeros :

Find *nearest* coefficient set \tilde{a}_{\min} with $p(\tilde{z}; \tilde{a}_{\min}) = 0$

$$\delta_{\min}(\tilde{z}) := \|\tilde{a}_{\min} - \bar{a}\|_e$$

Validity checking permits the following

Algorithmic Scheme

- compute **approximate result** \tilde{z} *somehow*
(approximate algorithm, floating-point)
- check validity of \tilde{z}

if **not valid** : refine \tilde{z} , go to check

else : **accept** \tilde{z} as **valid result**

Valid Result Sets

Problem with **empirical** data $(\bar{a}; e)$:

$$Z_\delta(\bar{a}; e) := \{ \tilde{z} \in \mathcal{Z} : \tilde{z} \text{ a } \delta\text{-valid result} \}$$

size of $Z_\delta(\bar{a}; e)$ indicates **condition** of result $F(\bar{a})$

“Pseudospectra” introduced for matrix eigenvalues in 90’s

Valid result sets for *distinct* results **not** meaningful when these results are to be used *jointly*.

Must be *jointly validated* if jointly used ($k \leq M$):

$$\mathcal{M}(\{\tilde{z}_1, \dots, \tilde{z}_k\}) := \mathcal{M}(\tilde{z}_1) \cap \dots \cap \mathcal{M}(\tilde{z}_k)$$

E.g. *Several* polynomial zeros \tilde{z}_κ , $\kappa = 1(1)k$:

$$\mathcal{M}(\{\tilde{z}_1, \dots, \tilde{z}_k\}) := \{\tilde{a} \in \mathcal{A} : p(\tilde{z}_\kappa; \tilde{a}) = 0, \kappa = 1(1)k\}$$

$\{\tilde{z}_1, \dots, \tilde{z}_k\}$ *simultaneously* δ -valid iff

$$N_\delta(\bar{a}, e) \cap \mathcal{M}(\{\tilde{z}_1, \dots, \tilde{z}_k\}) \neq \emptyset$$

$$Z_\delta[\{\tilde{z}_1, \dots, \tilde{z}_k\}] \not\subseteq Z_\delta[\tilde{z}_1] \cup \dots \cup Z_\delta[\tilde{z}_k]$$

M pol. zeros: $\mathcal{M}(\{\tilde{z}_1, \dots, \tilde{z}_M\}) = \{\tilde{a}^*\}$, interpolating pol.

Validity for Ill-Posed Problems

Type 1: $\mathcal{S} := \text{domain}(F) \subset \mathcal{A}$, $\dim \mathcal{S} =: d < M$

$$\mathcal{M}(\tilde{z}) := \{ \tilde{a} \in \mathcal{S} : F(\tilde{a}) = \tilde{z} \} \subset \mathcal{S} \subset \mathcal{A}$$

E.g., *Multiple* zero of $p(x; \bar{a})$: discriminant $(p, p') = 0$
now, refinement becomes (linearized) *optimization*

Type 2: $\bar{\mathcal{S}} := \text{image}(F) \subset \mathcal{Z}$, $\dim \bar{\mathcal{S}} =: \bar{d} < m$

$\tilde{z} \notin \bar{\mathcal{S}}$ is *not a result*, $\mathcal{M}(\tilde{z})$ is *not defined*

Must check validity of \tilde{z} in further use

Refine \tilde{z} towards $F(\bar{a})$ on $\bar{\mathcal{S}}$

E.g., numerically computed Groebner basis polynomials do *not* satisfy syzygies, may refine towards “exact” syzygies

Type 3: F discontinuous, e.g. discrete

\tilde{z} valid if $\mathcal{M}(\tilde{z}) \supset N_\delta(\bar{a}, \mathbf{e})$, $\delta = O(1)$

else: Check validity of “more degenerate” results

E.g., rank of **empirical** matrix = *lowest* valid rank

Basic Principles

interpret quantitative problem as a

data \rightarrow result map $F : \mathcal{A} \rightarrow \mathcal{Z}$

Design approximate map \tilde{F}

Evaluate \tilde{F} for specified data

obtain approximate result \tilde{z}

Compute backward error of \tilde{z}

accept or refine \tilde{z}

Note :

Want exact results (except for round-off)

for valid algebraic problems

i.e. for data very close to specified ones

Conclusions

Approximate Commutative Algebra
is meaningful and feasible
in the context of
quantitative problems
with coefficients with a natural topology
it is indispensable for
empirical problems